

A Model Proof in Logic \mathbb{E}

This is a proof of

$$\vdash (\neg p \equiv p) \equiv false$$

which is (3.15) where P is p . Typically, (3.1) and (3.1)' are used transparently. In this proof, detail concerning their use is explicitly given.

$$\begin{aligned}
 & (\neg p \equiv p) \equiv false \\
 = & \left\langle \begin{array}{l} \text{(3.8) and Leibniz where } E \text{ is } (\neg p \equiv p) \equiv r \\ \text{gives } \vdash ((\neg p \equiv p) \equiv false) \equiv ((\neg p \equiv p) \equiv \neg true). \end{array} \right\rangle \\
 & (\neg p \equiv p) \equiv \neg true \\
 = & \left\langle \begin{array}{l} \text{(3.2) where } P \text{ is } \neg p \text{ and } Q \text{ is } p \text{ and Leibniz where } E \text{ is } r \equiv \neg true \\ \text{gives } \vdash ((\neg p \equiv p) \equiv \neg true) \equiv ((p \equiv \neg p) \equiv \neg true). \end{array} \right\rangle \\
 & (p \equiv \neg p) \equiv \neg true \\
 = & \left\langle \text{(3.1) where } P \text{ is } p, Q \text{ is } \neg p, \text{ and } R \text{ is } \neg true. \right\rangle \\
 & p \equiv (\neg p \equiv \neg true) \\
 = & \left\langle \begin{array}{l} \text{(3.11) where } P \text{ is } p \text{ and } Q \text{ is } \neg true \text{ and Leibniz where } E \text{ is } p \equiv r \\ \text{gives } \vdash (p \equiv (\neg p \equiv \neg true)) \equiv (p \equiv (p \equiv \neg true)). \end{array} \right\rangle \\
 & p \equiv (p \equiv \neg \neg true) \\
 = & \left\langle \begin{array}{l} \text{(3.12) where } P \text{ is } true \text{ and Leibniz where } E \text{ is } p \equiv (p \equiv r) \\ \text{gives } \vdash (p \equiv (p \equiv \neg \neg true)) \equiv (p \equiv (p \equiv true)). \end{array} \right\rangle \\
 & p \equiv (p \equiv true) \\
 = & \left\langle \text{(3.1)' where } P \text{ is } p, Q \text{ is } p \text{ and } R \text{ is } true. \right\rangle \\
 & (p \equiv p) \equiv true \\
 = & \left\langle \begin{array}{l} \text{(3.3) where } P \text{ is } p \text{ and Leibniz where } E \text{ is } (p \equiv p) \equiv r \\ \text{gives } \vdash ((p \equiv p) \equiv true) \equiv ((p \equiv p) \equiv (p \equiv p)). \end{array} \right\rangle \\
 & (p \equiv p) \equiv (p \equiv p) .
 \end{aligned}$$

Transitivity applied four times gives

$$\vdash ((\neg p \equiv p) \equiv false) \equiv ((p \equiv p) \equiv (p \equiv p)) .$$

But (3.5) where P is $p \equiv p$ gives $\vdash (p \equiv p) \equiv (p \equiv p)$. By Equanimity,

$$\vdash \neg p \equiv p \equiv false .$$

Note that the use of Transitivity four times and Equanimity once, can be replaced by the use of Equanimity five times.

Note that (3.1)' is the Axiom Scheme

$$\vdash (P \equiv (Q \equiv R)) \equiv ((P \equiv Q) \equiv R) .$$