

Chapter 12 Exercise List SOLUTIONS

1. Let P be an expression of type Boolean. Suppose n dnof in P . Let

$$x_1, x_2, x_3, \dots, x_n, x_{n+1}, \dots$$

be an infinite sequence of variables that do not occur freely in P .

Define boolean expressions $E.j$ inductively (recursively) by defining

$$\begin{aligned} E.0 & \text{ to be } P, \text{ and} \\ E(n+1) & \text{ to be } (\forall x_{n+1} | : E.n). \end{aligned}$$

(Examples: $E.2$ could be $(\forall x | : (\forall y | : P))$;

$E.3$ could be $(\forall z | : (\forall x | : (\forall y | : P)))$, etc.)

- (a) Prove by Mathematical Induction that $\vdash (\forall n : \mathbb{N} | : P \Rightarrow E.n)$.

Answer: Note that $P.n$ is $P \Rightarrow E.n$ and that in particular, $P.0$ is $P \Rightarrow E.0$.

For the base case, we have

$$\begin{aligned} & P \Rightarrow E.0 \\ = & \langle \text{Recursive Definition} \rangle \\ & P \Rightarrow P \\ = & \langle (3.71) \rangle \\ & \text{true} \end{aligned}$$

For the Induction Step, assume $P.n$, i.e., $P \Rightarrow E.n$.

$$\begin{aligned} & P \Rightarrow E.(n+1) \\ = & \langle \text{Recursive Definition} \rangle \\ & P \Rightarrow (\forall x_{n+1} | : E.n) \\ = & \langle (3.59) \rangle \\ & \neg P \vee (\forall x_{n+1} | : E.n) \\ = & \langle (9.5) \rangle \\ & (\forall x_{n+1} | : \neg P \vee E.n) \\ = & \langle (3.59) \rangle \\ & (\forall x_{n+1} | : P \Rightarrow E.n) \end{aligned}$$

which is a theorem by Strong (9.16).

- (b) Give a complete, careful, syntactic proof that $\vdash P \Rightarrow E.143$.
(You may assume that 143 is a constant of type \mathbb{N} .)

Answer:

By (9.13), we have $\vdash (\forall n | P.n) \Rightarrow P.143$.

By Part (a), $\vdash (\forall n | : P.n)$.

By Modus Ponens, we have $\vdash P.143$, i.e., $\vdash P \Rightarrow E.143$.

2. For sets $S_i, i \in \mathbb{N}$, define $\bigcap_{i=0}^n S_i$ recursively by

$$\bigcap_{i=0}^0 S_i = S_0 \quad \text{and} \quad \bigcap_{i=0}^{n+1} S_i = \left(\bigcap_{i=0}^n S_i \right) \cap S_{n+1} .$$

Prove by Mathematical Induction that

$$\left(\forall n \mid : \bigcap_{i=0}^n (T_i \cap V) = \left(\bigcap_{i=0}^n T_i \right) \cap V \right) .$$

Answer: Base Case:

$$\vdash \left(\forall n \mid : \bigcap_{i=0}^0 (T_i \cap V) = \left(\bigcap_{i=0}^0 T_i \right) \cap V \right) .$$

$$\begin{aligned} & \bigcap_{i=0}^0 (T_i \cap V) = \left(\bigcap_{i=0}^0 T_i \right) \cap V \\ & = \langle \text{Definition} \rangle \\ & T_0 \cap V = T_0 \cap V . \end{aligned}$$

Induction Step:

$$\vdash \left(\forall j \mid 1 \leq j \leq n : \bigcap_{i=0}^j (T_i \cap V) = \left(\bigcap_{i=0}^j T_i \right) \cap V \right) \Rightarrow \bigcap_{i=0}^n (T_i \cap V) = \left(\bigcap_{i=0}^n T_i \right) \cap V .$$

Assume $\left(\forall j \mid 1 \leq j \leq n : \bigcap_{i=0}^j (T_i \cap V) = \left(\bigcap_{i=0}^j T_i \right) \cap V \right)$ from which $\bigcap_{i=0}^n (T_i \cap V) = \left(\bigcap_{i=0}^n T_i \right) \cap V$.

We are using = for **type** set in the left margin.

$$\begin{aligned} & \bigcap_{i=0}^{n+1} (T_i \cap V) \\ & = \langle \text{Recursive definition} \rangle \\ & \bigcap_{i=0}^n (T_i \cap V) \cap (T_{n+1} \cap V) \\ & = \langle \text{Assumption} \rangle \\ & \left(\bigcap_{i=0}^n T_i \right) \cap V \cap T_{n+1} \cap V \\ & = \langle (11.33), (11.34), (11.35) \rangle \\ & \left(\bigcap_{i=0}^n T_i \right) \cap T_{n+1} \cap V \\ & = \langle \text{Recursive definition} \rangle \\ & \left(\bigcap_{i=0}^{n+1} T_i \right) \cap V \end{aligned}$$

3. This exercise is from Rosen, **Discrete Mathematics and Its Applications**. Its solution is nontrivial.

$$\text{Prove, } \sum_{\emptyset \subset \{a_1, \dots, a_k\} \subseteq \{1, 2, \dots, n\}} \frac{1}{a_1 a_2 \cdots a_k} = n.$$

Take x, y, z to be fresh variables. For $y \subseteq \{1, \dots, n\}$ and $y \neq \emptyset$, define

$$\text{frac.}y = \frac{1}{(\prod x \mid x \in y : x)}.$$

Use Mathematical Induction to prove

$$\vdash (\forall n \mid n \geq 1 : (+y \mid y \subseteq \{1, \dots, n\} \wedge y \neq \emptyset : \text{frac.}y) = n).$$

Hint: Look at how one proceeds from Case $n = 2$ to Case $n = 3$. The nonempty subsets of $\{1, 2, 3\}$ are $\{1\}, \{2\}, \{1, 2\}$ which are the nonempty subsets of $\{1, 2\}$, as well as $\{3\}$, and $\{1, 3\}, \{2, 3\}, \{1, 2, 3\}$ which are those subsets of $\{1, 2, 3\}$ which are the union of nonempty subsets of $\{1, 2\}$ with $\{3\}$. If we know $\frac{1}{1} + \frac{1}{2} + \frac{1}{1 \cdot 2} = 2$, then

$$\begin{aligned} & \frac{1}{1} + \frac{1}{2} + \frac{1}{1 \cdot 2} + \frac{1}{3} + \frac{1}{1 \cdot 3} + \frac{1}{2 \cdot 3} + \frac{1}{1 \cdot 2 \cdot 3} \\ &= \left(\frac{1}{1} + \frac{1}{2} + \frac{1}{1 \cdot 2} \right) + \left(\frac{1}{3} \right) + \left(\frac{1}{1} + \frac{1}{2} + \frac{1}{1 \cdot 2} \right) \cdot \frac{1}{3} \\ &= 2 + \frac{1}{3} + 2 \cdot \frac{1}{3} = 2 + 1 = 3. \end{aligned}$$

The central step in the general case uses this change of dummy Lemma,

$$\begin{aligned} & \vdash (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} : \text{frac.}y) \\ &= (+z \mid z \subseteq \{1, \dots, n\} \wedge z \neq \emptyset : \frac{1}{n+1} \text{frac.}z). \end{aligned}$$

Answer: Take x, y, z to be fresh variables.

The proof is by Mathematical Induction. The base case is $n = 1$. Recall from discussion in class of $\mathcal{P}\{E\}$, that

$$\vdash y \subseteq \{1\} \wedge y \neq \emptyset \equiv y = \{1\}.$$

Now

$$\begin{aligned} & (+y \mid y \subseteq \{1\} \wedge y \neq \emptyset : \text{frac.}y) = 1 \\ &= \langle \vdash y \subseteq \{1\} \wedge y \neq \emptyset \equiv y = \{1\} \rangle \\ & \quad (+y \mid y = \{1\} : \text{frac.}y) = 1 \\ &= \langle (8.14) \rangle \\ & \quad \text{frac.}\{1\} = 1 \\ &= \langle \text{Definition of frac.}y \rangle \\ & \quad \frac{1}{(\prod x \mid x \in \{1\} : x)} = 1 \\ &= \langle \text{Lemma: } \vdash x \in \{1\} \equiv x = 1 \rangle \end{aligned}$$

$$\begin{aligned}
& \frac{1}{(\prod x \mid x = 1 : x)} = 1 \\
& = \langle (8.14) \rangle \\
& \frac{1}{1} = 1 .
\end{aligned}$$

For the induction step we prove

$$\vdash (+y \mid y \subseteq \{1, \dots, n\} \wedge y \neq \emptyset : \text{frac}.y) = n \Rightarrow (+y \mid y \subseteq \{1, \dots, n+1\} \wedge y \neq \emptyset : \text{frac}.y) = n+1 .$$

Assume $(+y \mid y \subseteq \{1, \dots, n\} \wedge y \neq \emptyset : \text{frac}.y) = n$. Leibniz (8.12) may be used for quantifications over y and z as neither of these variables appear free in the assumption.

Observe that under this assumption

$$\begin{aligned}
& (+y \mid y \subseteq \{1, \dots, n+1\} \wedge y \neq \emptyset : \text{frac}.y) = n+1 \\
& = \langle \vdash P \equiv (P \wedge \neg Q) \vee (P \wedge Q) \rangle \\
& \quad (+y \mid (y \subseteq \{1, \dots, n+1\} \wedge n+1 \notin y \wedge y \neq \emptyset) \\
& \quad \quad \vee (y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \emptyset) : \text{frac}.y) = n+1 \\
& = \langle (8.16) \rangle \\
& \quad (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \notin y \wedge y \neq \emptyset : \text{frac}.y) \\
& \quad \quad + (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \emptyset : \text{frac}.y) = n+1 \\
& = \langle \text{Lemma: } \vdash y \subseteq \{1, \dots, n+1\} \wedge n+1 \notin y \wedge y \neq \emptyset \equiv y \subseteq \{1, \dots, n\} \wedge y \neq \emptyset \rangle \\
& \quad (+y \mid y \subseteq \{1, \dots, n\} \wedge y \neq \emptyset : \text{frac}.y) \\
& \quad \quad + (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \emptyset : \text{frac}.y) = n+1 \\
& = \langle \text{Assumption} \rangle \\
& \quad n + (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \emptyset : \text{frac}.y) = n+1 \\
& = \langle \text{Algebra} \rangle \\
& \quad (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \emptyset : \text{frac}.y) = 1 .
\end{aligned}$$

We'll be done once we prove

$$\vdash (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \emptyset : \text{frac}.y) = 1 .$$

A routine proof gives

$$\vdash y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \emptyset \equiv y = \{n+1\} \vee (y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\}) .$$

Applying (8.16),

$$\begin{aligned}
& \vdash (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \emptyset : \text{frac}.y) \\
& = (+y \mid y = \{n+1\} : \text{frac}.y) + (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} : \text{frac}.y) .
\end{aligned}$$

Since by (8.14), $\vdash (+y \mid y = \{n+1\} : \text{frac}.y) = \frac{1}{n+1}$, we need only prove

$$\begin{aligned}
& \vdash (+y \mid y \subseteq \{1, \dots, n\} \wedge y \neq \emptyset : \text{frac}.y) = n \\
& \Rightarrow (+y \mid y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} : \text{frac}.y) = \frac{n}{n+1} .
\end{aligned}$$

The result follows easily from

$$\begin{aligned} & \vdash (+y | y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} : \text{frac}.y) \\ & = (+z | z \subseteq \{1, \dots, n\} \wedge z \neq \emptyset : \frac{1}{n+1} \text{frac}.z) . \end{aligned}$$

We have

$$\begin{aligned} & (+z | z \subseteq \{1, \dots, n\} \wedge z \neq \emptyset : \frac{1}{n+1} \text{frac}.z) \\ = & \langle (8.14) \rangle \\ & (+z | z \subseteq \{1, \dots, n\} \wedge z \neq \emptyset : (+y | y = z \cup \{n+1\} : \frac{1}{n+1} \text{frac}.z)) \\ = & \langle (8.20) \rangle \\ & (+z, y | z \subseteq \{1, \dots, n\} \wedge z \neq \emptyset \wedge y = z \cup \{n+1\} : \frac{1}{n+1} \text{frac}.z) \\ = & \left\langle \begin{array}{l} \text{Lemma: } \vdash z \subseteq \{1, \dots, n\} \wedge z \neq \emptyset \wedge y = z \cup \{n+1\} \\ \quad \equiv y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} \wedge z = y - \{n+1\} \end{array} \right\rangle \\ & (+z, y | y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} \wedge z = y - \{n+1\} : \frac{1}{n+1} \text{frac}.z) \\ = & \langle \vdash (\star x, y | R : P) = (\star y, x | R : P) \rangle \\ & (+y, z | y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} \wedge z = y - \{n+1\} : \frac{1}{n+1} \text{frac}.z) \\ = & \langle (8.20) \rangle \\ & (+y | y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} : (+z | z = y - \{n+1\} : \frac{1}{n+1} \text{frac}.z)) \\ = & \langle (8.14) \rangle \\ & (+y | y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} : \frac{1}{n+1} \text{frac}.(y - \{n+1\})) \\ = & \langle \text{Lemma: } \vdash n+1 \in y \wedge y \neq \{n+1\} \Rightarrow \frac{1}{n+1} \text{frac}.(y - \{n+1\}) = \text{frac}.y \rangle \\ & (+y | y \subseteq \{1, \dots, n+1\} \wedge n+1 \in y \wedge y \neq \{n+1\} : \text{frac}.y) \end{aligned}$$

Exercise 8, Section 1.2.1 of Donald Knuth's *The Art of Programming, Volume 1* reads

(a) Prove the following theorem of Nicomachus (C.E. c.100) by induction:

$$1^3 = 1, \quad 2^3 = 3 + 5, \quad 3^3 = 7 + 9 + 11, \quad 4^3 = 13 + 15 + 17 + 19, \quad \text{etc.}$$

(b) Use this result to prove the remarkable formula

$$1^3 + 2^3 + \dots + n^3 = (1 + 2 + \dots + n)^2 .$$

Recall from the 1090 Review Homework,

$$\vdash (+k | 1 \leq k \leq n : 2k - 1) = n^2 .$$

The following exercises provide a proof of (b) based on Knuth's approach. As in the previous problem, a change of dummy result lies at the core of the proof part (c).

4. (a) Use Mathematical Induction to prove that provided k d.n.o.f. in P ,

$$\vdash (+k | 1 \leq k \leq n : P) = n \cdot P .$$

Answer:

Case $n = 0$:

$$\begin{aligned} & (+k | 1 \leq k \leq 0 : P) = 0 \cdot P \\ = & \langle \text{Arithmetic, } \vdash 1 \leq k \leq 0 \equiv \text{false} \rangle \\ & (+k | \text{false} : P) = 0 \\ = & \langle (8.13) \rangle \\ & 0 = 0 . \end{aligned}$$

Assume $(\forall i | 0 \leq i \leq n : (+k | 1 \leq k \leq i : P) = i \cdot P)$ from which it follows that $(+k | 1 \leq k \leq n : P) = n \cdot P$.

$$\begin{aligned} & (+k | 1 \leq k \leq n + 1 : P) \\ = & \langle (8.23) \rangle \\ & (+k | 1 \leq k \leq n : P) + P[k := n + 1] \\ = & \langle \text{Assumption, } k \text{ d.n.o.f. in } P \rangle \\ & n \cdot P + P \\ = & \langle \text{Algebra} \rangle \\ & (n + 1) \cdot P \end{aligned}$$

- (b) Prove

$$\vdash (+j | 1 \leq j \leq n : j^3) = (+j | 1 \leq j \leq n : (+k | 1 \leq k \leq j : j^2 - j + (2k - 1))) .$$

Answer:

$$\begin{aligned} & (+k | 1 \leq k \leq j : j^2 - j + (2k - 1)) \\ = & \langle (8.15) \rangle \\ & (+k | 1 \leq k \leq j : j^2 - j) + (+k | 1 \leq k \leq j : 2k - 1) \\ = & \langle \text{Problem 5, 1090 Review} \rangle \\ & j \cdot (j^2 - j) + j^2 \\ = & \langle \text{Algebra} \rangle \\ & j^3 . \end{aligned}$$

The result follows using Leibniz.

(c) Prove

$$\vdash (n+1)^2 - (n+1) = 2 \cdot (+k \mid 1 \leq k \leq n : k)$$

and use Mathematical Induction to prove

$$\vdash (+j \mid 1 \leq j \leq n : (+k \mid 1 \leq k \leq j : j^2 - j + (2k-1))) = (+k \mid 1 \leq k \leq (+j \mid 1 \leq j \leq n : j) : 2k-1)$$

Answer:

$$\begin{aligned} & (n+1)^2 - (n+1) \\ = & \langle \text{Quoted result from 1090 Review} \rangle \\ & (+k \mid 1 \leq k \leq n+1 : 2k-1) - (n+1) \\ = & \langle (8.23) \rangle \\ & (+k \mid 1 \leq k \leq n : 2k-1) + (2(n+1) - 1) - (n+1) \\ = & \langle (8.15), \text{Algebra} \rangle \\ & (+k \mid 1 \leq k \leq n : 2k) + (+k \mid 1 \leq k \leq n : -1) + n \\ = & \langle \text{Problem 5} \rangle \\ & (+k \mid 1 \leq k \leq n : 2k) - n + n \\ = & \langle \text{Algebra} \rangle \\ & (+k \mid 1 \leq k \leq n : k+k) \\ = & \langle (8.15), \text{Algebra} \rangle \\ & 2 \cdot (+k \mid 1 \leq k \leq n : k) . \end{aligned}$$

To prove

$$\vdash (+j \mid 1 \leq j \leq n : (+k \mid 1 \leq k \leq j : j^2 - j + (2k-1))) = (+k \mid 1 \leq k \leq (+j \mid 1 \leq j \leq n : j) : 2k-1),$$

consider the Base Case, $n = 0$.

$$\begin{aligned} & (+j \mid 1 \leq j \leq 0 : (+k \mid 1 \leq k \leq j : j^2 - j + (2k-1))) \\ = & (+k \mid 1 \leq k \leq (+j \mid 1 \leq j \leq 0 : j) : 2k-1) \\ = & \langle \text{Arithmetic, } \vdash 1 \leq j \leq 0 \equiv \text{false} \rangle \\ & (+j \mid \text{false} : (+k \mid 1 \leq k \leq j : j^2 - j + (2k-1))) \\ = & (+k \mid 1 \leq k \leq (+j \mid \text{false} : j) : 2k-1) \\ = & \langle (8.13) \rangle \\ & 0 = (+k \mid 1 \leq k \leq 0 : 2k-1) \\ = & \langle \text{Arithmetic, } \vdash 1 \leq k \leq 0 \equiv \text{false} \rangle \\ & 0 = (+k \mid \text{false} : 2k-1) . \end{aligned}$$

Now assume

$$\begin{aligned} & (\forall i \mid 0 \leq i \leq n : (+j \mid 1 \leq j \leq i : (+k \mid 1 \leq k \leq j : j^2 - j + (2k-1)))) \\ & = (+k \mid 1 \leq k \leq (+j \mid 1 \leq j \leq i : j) : 2k-1) , \end{aligned}$$

from which it follow that

$$\begin{aligned}
(+j|1 \leq j \leq n : (+k|1 \leq k \leq j : j^2 - j + (2k-1))) &= (+k|1 \leq k \leq (+j|1 \leq j \leq n : j) : 2k-1) . \\
&= (+j|1 \leq j \leq n+1 : (+k|1 \leq k \leq j : j^2 - j + (2k-1))) \\
&= \langle (8.23) \rangle \\
&= (+j|1 \leq j \leq n : (+k|1 \leq k \leq j : j^2 - j + (2k-1))) \\
&\quad + (+k|1 \leq k \leq n+1 : (n+1)^2 - (n+1) + (2k-1)) \\
&= \langle \text{previous result} \rangle \\
&= (+k|1 \leq k \leq (+j|1 \leq j \leq n : j) : 2k-1) \\
&\quad + (+k|1 \leq k \leq n+1 : 2 \cdot (+j|1 \leq j \leq n : j) + (2k-1)) \\
&= \langle \text{Algebra} \rangle \\
&= (+k|1 \leq k \leq (+j|1 \leq j \leq n : j) : 2k-1) \\
&\quad + (+k|1 \leq k \leq n+1 : 2 \cdot ((+j|1 \leq j \leq n : j) + k) - 1) \\
&= \langle \text{Change of Dummy, proof similar to Problem 4(a), 1090 Review} \rangle \\
&= (+k|1 \leq k \leq (+j|1 \leq j \leq n : j) : 2k-1) \\
&\quad + (+k|(+j|1 \leq j \leq n : j) + 1 \leq k \leq (+j|1 \leq j \leq n : j) + (n+1) : 2k-1) \\
&= \langle (8.23) \rangle \\
&= (+k|1 \leq k \leq (+j|1 \leq j \leq n : j) : 2k-1) \\
&\quad + (+k|(+j|1 \leq j \leq n : j) + 1 \leq k \leq (+j|1 \leq j \leq n+1 : j) : 2k-1) \\
&= \langle (8.16) \rangle \\
&= (+k|1 \leq k \leq (+j|1 \leq j \leq n+1 : j) : 2k-1)
\end{aligned}$$

(d) Prove

$$\vdash (+j|1 \leq j \leq n : j^3) = (+j|1 \leq j \leq n : j)^2 .$$

Answer: Problems 2. and 3. give

$$\vdash (+j|1 \leq j \leq n : j^3) = (+k|1 \leq k \leq (+j|1 \leq j \leq n+1 : j) : 2k-1) .$$

But by the result from Homework 2,

$$\vdash (+k|1 \leq k \leq (+j|1 \leq j \leq n+1 : j) : 2k-1) = (+j|1 \leq j \leq n+1 : j)^2 .$$

By (1.4),

$$\vdash (+j|1 \leq j \leq n : j^3) = (+j|1 \leq j \leq n : j)^2 .$$