Quantale-valued Topological Spaces via Closure and Convergence

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Abstract

For a quantale \( V \) we introduce \( V \)-valued topological spaces via \( V \)-valued point-set-distance functions and, when \( V \) is completely distributive, characterize them in terms of both, so-called closure towers and ultrafilter convergence relations. When \( V \) is the two-element chain \( 2 \), the extended real half-line \([0, \infty]\), or the quantale \( \Delta \) of distance distribution functions, the general setting produces known and new results on topological spaces, approach spaces, and the only recently considered probabilistic approach spaces, as well as on their functorial interactions with each other.

Keywords: quantale, \( V \)-valued closure space, \( V \)-valued topological space, discrete \( V \)-presheaf monad, lax distributive law, lax \((\lambda, V)\)-algebra, probabilistic approach space, algebraic functor, change-of-base functor.

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1. Introduction

Lowen’s [17] approach spaces provide an ideal synthesis of Lawvere’s [16] presentation of metric spaces (as small \([0, \infty]\)-enriched categories) and the Manes-Barr [19, 1] presentation of topological spaces in terms of ultrafilter convergence, as demonstrated first in [4]; see also [12]. Several authors have investigated probabilistic generalizations of these concepts (see in particular [20, 3, 11, 14]), which suggests that a quantale-based study of generalized topological spaces should be developed, in order to treat these and other new concepts efficiently in a unified manner, in terms of both, “distance” or “closure”, and “convergence”. In this paper we provide such a treatment, working with an arbitrary quantale \( V = (V, \otimes, k) \) which, for the main results of the paper, is required to be completely distributive. For \( V = 2 \) the two-element chain, our results reproduce the equivalence of the descriptions of topologies in terms of closure and ultrafilter convergence; for \( V = [0, \infty] \) (ordered by the natural \( \geq \) and structured by \( + \) as the quantalic \( \delta \)), one obtains the known equivalent descriptions of approach spaces in terms of point-set distances and of ultrafilter convergence; for \( V = \Delta \) the quantale of distance distribution functions \( \varphi : [0, \infty] \to [0, 1] \), required to satisfy the left-continuity condition \( \varphi(\beta) = \sup_{\alpha < \beta} \varphi(\alpha) \) for all \( \beta \in [0, \infty] \), the corresponding equivalence is established here also for probabilistic approach spaces. A major advantage of working in the harmonized context of a general quantale is that it actually makes the proofs more transparent to us than if they were carried out in the concrete quantales that we are interested in.

While this paper is built on the methods of monoidal topology as developed in [6, 5, 12] and elsewhere (see in particular [13]), in this paper we emphasize the lax-algebraic setting presented in [22], which is summarized in this paper to the extent needed. This setting is in fact well motivated by Lowen’s original axioms for an approach space \((X, \delta)\) in terms of its point-set distance function \( \delta : X \times PX \to [0, \infty] \), listed in [17] with \( PX = 2^X \), as follows:

\begin{align*}
(D1) \quad \forall x \in X : \quad \delta(x, \{x\}) &= 0, \\
(D2) \quad \forall x \in X : \quad \delta(x, \emptyset) &= \infty,
\end{align*}

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(D3) \( \forall x \in X, A, B \subseteq X : \delta(x, A \cup B) = \min(\delta(x, A), \delta(x, B)) \).
(D4) \( \forall x \in X, A \subseteq X, \varepsilon \in [0, \infty] : \delta(x, A) \leq \delta(x, A^{(\varepsilon)}) + \varepsilon \), where \( A^{(\varepsilon)} := \{ x \in X | \delta(x, A) \leq \varepsilon \} \).

Since (D2), (D3) require \( \delta(x, -) : (PX, \subseteq) \rightarrow ([0, \infty], \geq) \) to preserve finite joins for every fixed \( x \in X \), we are led to describe \( \delta \) equivalently as a function

\[
c : PX \rightarrow [0, \infty]^X
\]

which also avoids the quantification over \( x \) in each of these axioms. Now (D1) and (D4) may be interpreted as the reflexivity and transitivity axioms for a lax \((\mathbb{P}, [0, \infty])\)-algebra in the sense of [12], where \( \mathbb{P} \) is the powerset monad of \text{Set}, suitably extended to \([0, \infty]\)-valued relations of sets. Equivalently, as we will show in this paper, (D1) and (D4) provide \( X \) with a \([0, \infty]\)-indexed closure tower (named so after the terminology used in [3, 24]), the members of which are collectively extensive, monotone and idempotent, in a sense that we make precise in the general context of a quantale in Proposition 2.6. In this way we obtain new characterizations of approach spaces and of probabilistic approach spaces in terms of closure, which we summarize at the end of Section 2.

Lowen [17] also gave the equivalent description of the structure of an approach space \( X \) in terms of a limit operator \( F_X \rightarrow (0, \infty)^X \), which assigns to every filter on \( X \) a function that provides for every \( x \in X \) a measure of "how far away \( x \) is from being a limit point" of the given filter. As first shown in [4], it suffices to restrict this operator to ultrafilters, so that the structure may in fact be given by a map

\[
\ell : UX \rightarrow [0, \infty]^X
\]
satisfying two axioms that correspond to the reflexivity and transitivity conditions for a lax \(((0, \infty), [0, \infty])\)-algebra structure on \( X \) as described in [12], with \( U \) denoting the ultrafilter monad of \text{Set}, understood to be laxly extended from maps to \([0, \infty]\)-valued relations.

The presentations \((*)\), \((**)*\) motivated the study of lax \((\lambda, V)\)-algebras in [22], i.e., of sets provided with a map

\[
c : TX \rightarrow V^X
\]
satisfying two basic axioms. Here, for a \text{Set}-monad \( T = (T, m, e) \) and the given quantale \( V \), \( \lambda \) is a lax distributive law of \( T \) over \( V \), which links \( T \) with \( V \), as encoded by the \( V \)-powerset monad \( P_V = (P_V, s, y) \), with \( P_V X = V^X \).

For \( T = P = P_2 \) and a naturally chosen lax distributive law, the corresponding lax \( V \)-algebras are \( V \)-valued closure spaces, satisfying the \( V \)-versions of (D1), (D4); we call them \( V \)-valued topological spaces\(^2\) when they also satisfy the \( V \)-versions of (D2), (D3). The main result of the paper (Theorem 3.6) describes them equivalently as the lax \( V \)-algebras with respect to a naturally chosen lax distributive law of the ultrafilter monad \( U \) over \( P_V \), provided that \( V \) is completely distributive. The relevant isomorphism of categories comes about as the restriction of an adjunction, the left-adjoint functor of which is an algebraic functor as discussed in [22] (in generalization of the well-known concept presented in [5, 12]). For \( V = \Delta \) our general result produces a new characterization of probabilistic approach spaces in terms of ultrafilter convergence (Corollary 3.7).

In the last section we study so-called change-of-base functors (see [5, 12, 22]) for the categories at issue in this paper. An application of our general result (Theorem 4.4) gives a unified proof for the known facts that \( \text{Top} \) may be fully embedded into \text{App} as a simultaneously reflective and coreflective subcategory which, in turn is reflectively and coreflectively embedded into \text{ProbApp}.

2. \( V \)-valued topological spaces via closure

Throughout the paper, let \( V = (V, \otimes, k) \) be a (unital but not necessarily commutative) quantale, i.e., a complete lattice with a monoid structure whose binary operation \( \otimes \) preserves suprema in each variable. There are no additional provisions for the tensor-neutral element \( k \) vis-à-vis the bottom and top elements in \( V \), i.e., we exclude neither the case \( k = \bot \) (so that \( |V| = 1 \)), nor \( k < T \). The \( V \)-powerset functor \( P_V : \text{Set} \rightarrow \text{Set} \) is given by

\[
(f : X \rightarrow Y) \mapsto (f_! : V^X \rightarrow V^Y), \quad f_!(\sigma)(y) = \bigvee_{x \in f^{-1}y} \sigma(x),
\]

\(^2\)We had preferred to omit the adjective "valued" but feel compelled not to do so in order to avoid potential confusion by some readers with the same term used in fuzzy topology for a different concept.
for all $\sigma : X \rightarrow V$, $y \in Y$. The functor $P_V$ carries a monad structure, given by

$$y_X : X \rightarrow V^X, \quad (y_X(y))(y) = \begin{cases} k & \text{if } y = x \\ \bot & \text{otherwise} \end{cases},$$

$$s_X : V^{V^X} \rightarrow V^X, \quad (s_X \Sigma)(x) = \bigvee_{\sigma \in V^X} \Sigma(\sigma) \otimes \sigma(x),$$

for all $x, y \in X$ and $\Sigma : V^X \rightarrow V$.

Let $\mathbb{T} = (T, m, e)$ be any monad on $\mathbf{Set}$. A lax distributive law $\lambda$ of $\mathbb{T}$ over $P_V = (P_V, s, y)$ (see [12, 22], and [2] for its original name giver) is a family of maps $\lambda_X : T(V^X) \rightarrow V^T (X \in \mathbf{Set})$ which, when one orders maps to a power of $V$ pointwise by the order of $V$, must satisfy the following conditions:

(a) $\forall f : X \rightarrow Y : \quad (Tf) \cdot \lambda_X \leq \lambda_Y \cdot T(f)$ (lax naturality of $\lambda$),

(b) $\forall X : \quad y_{TX} \leq \lambda_X \cdot Ty_X$ (lax $P_V$-unit law),

(c) $\forall X : \quad s_{TX} \cdot (\lambda_X) \cdot \lambda_Y \leq \lambda_X \cdot Ts_X$ (lax $P_V$-multiplication law),

(d) $\forall X : \quad (\lambda_Y) : (\lambda_X) \cdot \lambda_Y \leq \lambda_X \cdot e_{P_V}$ (lax $T$-unit law),

(e) $\forall X : \quad (m_X) \cdot \lambda_{TX} \cdot \lambda_Y \leq \lambda_X \cdot m_{P_V}$ (lax $T$-multiplication law),

(f) $\forall g, h : Z \rightarrow V^X : \quad g \leq h \Rightarrow \lambda_X \cdot Tg \leq \lambda_X \cdot Th$ (monotonicity).

\textbf{Remark 2.1.} Although we will make use of it only in the next section, let us mention here the fact that lax distributive laws of a $\mathbf{Set}$-monad $\mathbb{T} = (T, m, e)$ over $P_V$ correspond bijectively to lax extensions $\hat{T}$ of $\mathbb{T}$ to the category $V$-$\mathbf{Rel}$ of sets with $V$-valued relations $r : X \rightarrow Y$ as morphisms, which are equivalently displayed as maps $\hat{T} : Y \rightarrow P_V X$ (see [22] and Exercise III.1.1 in [12]). Given $\lambda$, the lax functor $\hat{T} : V$-$\mathbf{Rel} \rightarrow V$-$\mathbf{Rel}$ assigns to $r$ the $V$-relation $\hat{T}r : TX \rightarrow TY$ defined by

$$\hat{T}r = \lambda_X \cdot T \hat{r}.$$

Conversely, the lax distributive law $\lambda$ associated with $\hat{T}$ is given by

$$\lambda_X = \hat{r}_X \cdot T \hat{e}_X,$$

with $\hat{e}_X : X \rightarrow P_V X$ the evaluation $V$-relation: $\hat{e}_X(x, \sigma) = \sigma(x)$.

\textbf{Proposition 2.2.} The ordinary powerset monad $P = P_2$ distributes laxly over the $V$-powerset monad $P_V$, via

$$\alpha_X : P(V^X) \rightarrow V^{P^X}, \quad (\alpha_XS)(A) = \bigvee_{x \in A} \bigvee_{\sigma \in S} \sigma(x) \quad (S \subseteq V^X, A \subseteq X).$$

\textbf{Proof.} (a) For all $S \subseteq V^X, B \subseteq X$ one has

$$((Pf) \cdot \alpha_X(S))(B) = \bigvee_{A \subseteq X, f(A) = B} (\alpha_XS)(A) = \bigvee_{A \subseteq X, f(A) = B} \bigvee_{x \in A} \bigvee_{\sigma \in S} \sigma(x).$$

Lax naturality of $\alpha$ follows since, for every $A \subseteq X$ with $f(A) = B$,

$$\bigvee_{x \in A} \bigvee_{\sigma \in S} \sigma(x) \leq \bigvee_{y \in B} \bigvee_{\sigma \in S} \bigvee_{x \in f^{-1}y} \sigma(x) = \bigvee_{y \in B} \bigvee_{\sigma \in S} (f_\sigma)(y) = \alpha_Y(f_\sigma)(S)(B) = (\alpha_Y \cdot P(f))(S)(B).$$

(b) For all $A, B \subseteq X$,

$$\alpha_X \cdot P y_X(B)(A) = \bigvee_{x \in A} \bigvee_{y \in B} (y_X y)(x) = \begin{cases} \top & \text{if } A = \emptyset \\ k & \text{if } \emptyset \neq A \subseteq B \\ \bot & \text{otherwise} \end{cases} \supseteq (y_{P^X}B)(A).$$
(c) For all \( S \subseteq V^X, A \subseteq X \),
\[
(s_{PX} \cdot (\alpha_X) \cdot \alpha_{V^X}(S))(A) = \bigvee_{\tau \in V^X} ((\alpha_X)_! \cdot \alpha_{V^X}(S))(\tau) \otimes \tau(A)
\]
\[
= \bigvee_{\tau \in V^X} \bigg( \bigvee_{S \subseteq V^X, \alpha_X(S) = \tau} \alpha_{V^X}(S) \otimes \tau(A) \bigg)
\]
\[
\leq \bigvee_{S \subseteq V^X} \alpha_{V^X}(S) \otimes \alpha_X(S)(A)
\]
\[
= \bigg( \bigvee_{\tau \in V^X} \bigg( \bigg( \bigvee_{S \subseteq V^X} \bigvee_{\tau \in S} \Sigma(\tau) \otimes \bigvee_{x \in A} \sigma(x) \bigg) \bigg) \bigg) \subseteq A \subseteq X
\]
\[
\leq \bigvee_{S \subseteq V^X} \bigg( \bigvee_{x \in A} \bigvee_{\sigma \in S} \Sigma(\sigma) \otimes \sigma(x) \bigg)
\]
\[
= (\alpha_X \cdot s_X(S))(A).
\]

(d) With \( e_X : X \to PX \) denoting the map \( x \mapsto \{ x \} \), for all \( \sigma \in V^X, A \subseteq X \) one has
\[
(e_X)(\sigma)(A) = \bigvee_{x \subseteq X, \{ x \} = A} \sigma(x) = \begin{cases} \sigma(x) & \text{if } (\exists x : A = \{ x \}) \\ \perp & \text{otherwise} \end{cases} \leq \bigwedge_{x \in A} \sigma(x) = (\alpha_X \cdot e_{V^X}(\sigma))(A).
\]

(e) With \( m_X : PPX \to PX \) denoting the map \( \mathcal{A} \mapsto \bigcup \mathcal{A} \), for all \( S \subseteq V^X, A \subseteq X \) one has
\[
((m_X)_! \cdot \alpha_{PX} \cdot \mathcal{P}\alpha_X(S))(A) = \bigg( \bigvee_{B \subseteq \bigcup \mathcal{A}} \bigg( \bigvee_{B \subseteq \mathcal{A}} \bigg( \bigvee_{B \subseteq \mathcal{A}} \alpha_X(B) \bigg) \bigg) \bigg) \subseteq A \subseteq X
\]
\[
= \bigvee_{B \subseteq \bigcup \mathcal{A}} \bigg( \bigvee_{B \subseteq \mathcal{A}} \bigg( \bigvee_{B \subseteq \mathcal{A}} \bigvee_{S \subseteq V^X} \sigma(y) \bigg) \bigg) \subseteq A \subseteq X
\]
\[
= \bigvee_{S \subseteq V^X} \bigg( \bigvee_{x \in A} \bigvee_{B \subseteq \mathcal{A}} \bigg( \bigg( \bigvee_{x \in B} \bigvee_{\sigma \in S} \sigma(y) \bigg) \bigg) \bigg) \subseteq A \subseteq X
\]
\[
= (\alpha_X \cdot m_{V^X}(S))(A).
\]

But whenever \( x \in A = \bigcup \mathcal{A} \), so that \( x \in B_0 \) for some \( B_0 \in \mathcal{A} \), we have
\[
\bigvee_{B \subseteq \mathcal{A}} \bigg( \bigvee_{B \subseteq \mathcal{A}} \bigg( \bigvee_{B \subseteq \mathcal{A}} \bigvee_{S \subseteq V^X} \sigma(y) \bigg) \bigg) \leq \bigvee_{S \subseteq V^X} \bigg( \bigvee_{x \in A} \bigvee_{B \subseteq \mathcal{A}} \bigg( \bigg( \bigvee_{x \in B} \bigvee_{\sigma \in S} \sigma(y) \bigg) \bigg) \bigg)
\]
\[
\text{and may conclude } ((m_X)_! \cdot \alpha_{PX} \cdot \mathcal{P}\alpha_X(S))(A) \leq (\alpha_X \cdot m_{V^X}(S))(A).
\]

(f) From \( g \leq h \), for all \( C \subseteq Z, A \subseteq X \) one obtains immediately
\[
(\alpha_X(g))(A) = \bigvee_{x \in A} (g_x)(x) \leq \bigvee_{x \in A} (h_x)(x) = (\alpha_X(h))(A).
\]

\[\square\]

**Definition 2.3.** (22) Let \( \mathcal{A} \) be a lax distributive law of a **Set**-monad \( T \) over \( F_V \). A **lax** \((\mathcal{A}, V)\)-algebra \((X, c)\) is a set \( X \) with a map \( c : TX \to V^X \) satisfying
A lax homomorphism \( f : (X, c) \rightarrow (Y, d) \) of lax \( (\lambda, V) \)-algebras is a map \( f : X \rightarrow Y \) satisfying

\[
\text{(M)} \quad f \cdot c \leq d \cdot Tf
\]

(lax homomorphism law, monotonicity).

The resulting category is denoted by

\[
(\lambda, V)\text{-Alg}
\]

(2) A \( V \)-valued closure space \( (X, c) \) is a lax \( (\alpha, V) \)-algebra, with \( \alpha \) as in Proposition 2.2; it is a \( V \)-valued topological space if, in addition, \( c : PX \rightarrow V^X \) preserves finite joins:

\[
\forall x \in X, A, B \subseteq X : (c\emptyset)(x) = \bot \text{ and } c(A \cup B)(x) = (cA)(x) \lor (cB)(x).
\]

A lax \( \alpha \)-homomorphism of \( V \)-valued closure spaces is also called a contractive map. We obtain the category

\[
V\text{-Cls} = (\alpha, V)\text{-Alg}
\]

and its full subcategory by \( V\text{-Top} \).

**Remark 2.4.** If the lax distributive law \( \lambda \) is equivalently described as a lax extension \( \tilde{T} \) of \( T \) (see Remark 2.1), then

\[
(\lambda, V)\text{-Alg} \cong (\tilde{T}, V, \tilde{T})\text{-Cat}
\]

is the category of \( (\tilde{T}, V) \)-categories, as defined in [12]. Under this isomorphism (see [22], Prop. 6.8), the \( (\lambda, V) \)-structure \( c : TX \rightarrow PX \) corresponds to the \( V \)-relation \( a : TX \rightarrow X \) with \( a^\circ = c \) (where \( a^\circ : X \rightarrow TX \) is the converse of \( a \)), and (R) and (T) now read as

\[
k \leq a(e_X(x), x) \quad \text{and} \quad \tilde{T}a(\tilde{x}, \eta) \otimes a(\eta, z) \leq a(m_X \tilde{x}, z),
\]

for all \( \tilde{x} \in TTX, \eta \in TX, z \in X \).

The lax extension \( \tilde{T} : V\text{-Rel} \rightarrow V\text{-Rel} \) corresponding to \( \alpha \) of Proposition 2.2 is, after an easy computation, described by

\[
\tilde{T}r(A, B) = \bigwedge_{y \in B} \bigvee_{x \in A} r(x, y),
\]

for all \( V \)-relations \( r : X \rightarrow Y, A, B \subseteq X \). Consequently, stated elementwise, conditions (R), (T), (M) read for \( \lambda = \alpha \) as (R'), (T'), (M') of the following proposition. Stated equivalently as (R''), (T''), reflexivity and transitivity are interpreted as \( c \) being extensive and idempotent:

**Lemma 2.5.** A map \( c : PX \rightarrow V^X \) makes \( X \) a \( V \)-valued closure space if and only if, \( c \) satisfies

\[
\text{(R')} \quad \forall x \in X : \quad k \leq c(\{x\})(x), \quad \text{and}
\]

\[
\text{(T')} \quad \forall A \subseteq PX, B \subseteq X, z \in X : \quad \left( \bigwedge_{y \in B} \bigvee_{A \subseteq \{y\}} (cA)(y) \right) \otimes (cB)(z) \leq c(\cup A)(z)
\]

or, equivalently, if \( c \) satisfies

\[
\text{(R'')} \quad \forall A \subseteq X, x \in A : \quad k \leq (cA)(x), \quad \text{and}
\]

\[
\text{(T'')} \quad \forall A \subseteq X, B \subseteq X, x \in X : \quad \left( \bigwedge_{y \in B} (cA)(y) \right) \otimes (cB)(x) \leq (cA)(x).
\]

A map \( f : X \rightarrow Y \) of \( V \)-valued closure spaces \( (X, c), (Y, d) \) is contractive if, and only if,

\[
\text{(M')} \quad \forall x \in X, A \subseteq X : \quad (cA)(x) \leq (df)(A)(fx).
\]

**Proof.** (R')&(T')\implies(R'')&(T''): For (R''), considering \( x \in A \) and putting \( A := \{y \mid y \in A\}, B := \{x\} \) one obtains

\[
k = k \otimes k \leq (cA)(x) \text{ from (R'), (T').}
\]

For (T''), considering \( A, B \subseteq X \) and putting \( A = [A] \), one obtains (T'') from (T') immediately.
(R"")&(T"")⇒(R')&(T'): Trivially, (R") implies (R'). To check (T'), first notice that the hypothesis makes c monotone: whenever $B \subseteq A \subseteq X$, one has

$$cB(x) \leq k \otimes cB(x) \leq \left( \bigwedge_{y \in B} cA(y) \right) \otimes cB(x) \leq cA(x)$$

for all $x \in X$. Consequently, for all $A \subseteq PX$ and $z \in X$ one obtains

$$\left( \bigwedge_{y \in B \in A} cA(y) \right) \otimes cB(z) \leq \left( \bigwedge_{y \in B} c(A \cup B)(y) \right) \otimes cB(z) \leq c(A \cup B)(z).$$

\[\square\]

We can now describe the structure of $V$-valued closure spaces in terms of $V$-indexed closure towers\(^3\), as follows.

**Proposition 2.6.** (1) For a $V$-valued closure space $(X, c)$, with

$$c^A := \{ x \in X \mid (cA)(x) \geq v \} \quad (v \in V, A \subseteq X)$$

one obtains a family of maps $(c^v : PX \to PX)_{v \in V}$ satisfying

(C0) if $B \subseteq A$, then $c^v B \subseteq c^A$,

(C1) if $v \leq \bigvee_{i \in I} u_i$, then $\bigwedge_{i \in I} c^{u_i} A \subseteq c^A$,

(C2) $A \subseteq cA$,

(C3) $c^v c^A \subseteq c^{v \circ u} A$,

for all $A, B \subseteq X$ and $v, u, u_i \in V$ ($i \in I$).

(2) Conversely, for any family maps $(c^v : PX \to PX)_{v \in V}$ satisfying the conditions (C0)–(C3), putting

$$(cA)(x) := \bigvee \{ v \in V \mid x \in c^A \} \quad (A \subseteq X, x \in X)$$

makes $(X, c)$ a $V$-valued closure space.

(3) The correspondences of (1), (2) are inverse to each other. Under this bijection, contractivity of a map $f : X \to Y$ is equivalently described by the continuity condition

$$\forall A \subseteq X, v \in V : \ f(c^A) \subseteq d^\circ f(A).$$

**Proof.** (1) Clearly, (C2) follows from (R""), and for (C0) see the proof of Lemma 2.5. (C1) follows trivially from the definition of the closure tower, and for (C3) one puts $B := c^v A$ to obtain, with (T""), $v \otimes c(c^v A)(x) \leq (cA)(x)$ for all $x \in X$. Hence, for $x \in c^v(c^v A)$ one may conclude $v \otimes u \leq (cA)(x)$, which means $x \in c^{v \circ u} A$.

(2) (R") follows trivially from (C2). In order to show (T""), putting

$$v := \bigvee \{ v \in V \mid y \in c^A \},$$

for every $y \in B$ we obtain from (C1) $y \in c^v A$, and then, with $\tilde{v} := \bigwedge_{y \in B} v_y$ and (C0), $y \in c_{\tilde{v}} A \subseteq c^\tilde{v} A$, so that $B \subseteq c^\tilde{v} A$. Now, for every $u \in V$, (C0) and (C3) give $c^\tilde{v}(B) \subseteq c^\tilde{v}(c^A) \subseteq c^{v \circ u} A$. Consequently, for all $x \in c^\tilde{v} B$, we obtain $\tilde{v} \otimes u \leq cA(x)$ and conclude

$$\left( \bigwedge_{y \in B} \bigvee \{ v \in V \mid y \in c^A \} \right) \otimes \bigvee \{ u \in V \mid x \in c^\tilde{v} B \} = \tilde{v} \otimes \bigvee \{ u \in V \mid x \in c^\tilde{v} B \} \leq cA(x),$$

as desired.

(3) Given a $V$-valued-closure-space structure $c$ on $X$, let $(c^v)_{v \in V}$ be the closure tower as in (1) and denote by $\tilde{c}$ the structure obtained from that tower as in (2). Since trivially $x \in c^A(x \cup A)$, one easily concludes $(\tilde{c}A)(x) = (cA)(x)$ for

\(^3\)Our conditions differ from the ones used by Lowen [18] for his closure towers in the case $V = [0, \infty]$. 
all $A \subseteq X, x \in X$. Conversely, starting with a closure tower $(c^v)_{v \in V}$, forming the corresponding $V$-valued-closure-space structure $c$ as in (2) and then its induced closure tower $(\tilde{c}^v)_{v \in V}$ as in (1), we conclude

$$\tilde{c}^v A = \{x \in X \mid \bigvee \{u \in V \mid x \in c^u A \} \geq v \} \subseteq c^v A$$

for all $A \subseteq X$ from (C1), with the reverse inclusion holding trivially.

Finally, that $(M')$ implies the given continuity condition follows directly from the definitions. In turn, the continuity condition implies $(M')$ when being exploited for $v := (cA)(x)$, since then $x \in c^v A$ and therefore $fx \in d'(f(A))$, which means precisely $(M')$.

\textbf{Remark 2.7.} (1) Note that, for a $V$-valued closure space $(X, c)$, one has $c^\bot A = X$ for all $A \subseteq X$ (including $A = \emptyset$). Hence, in Proposition 2.6(2), it suffices to require (C0)–(C3) for all those $u, v, u_i \in V (i \in I)$ that are greater than $\bot$.

(2) If $c$ and $(c^v)_{v \in V}$ correspond to each other as in Proposition 2.6(1),(2), then (C3) may be written equivalently as

$$(C3') \quad v \otimes c(c^v A)(x) \leq (cA)(x),$$

for all $x \in X, A \subseteq X, v \in V$. Indeed, from (C3) one obtains

$$v \otimes c(c^v A)(x) = v \otimes \bigvee \{u \in V \mid x \in c^u A \} \leq \bigvee \{v \otimes u \mid u \in V, x \in c^u A \} \leq \bigvee \{w \in V \mid x \in c^w A \} = (cA)(x);$$

conversely, given $(C3')$, one has

$$x \in c^v(c^v A) \implies (cA)(x) \geq v \otimes c(c^v A)(x) \geq v \otimes u \implies x \in c^u A.$$

(3) For Lawvere’s quantale $[0, \infty]$, ordered by the natural $\geq$ and provided with $\otimes = +$, naturally extended to $\infty$, writing $\delta(x, A) = (cA)(x)$ one sees that condition $(C3')$ coincides with (D4) (see Introduction).

We are now ready to describe $V$-valued topological spaces in terms of closure towers, provided that $V$ is \textit{constructively completely distributive (ccd)}. Recall that the complete lattice $V$ is ccd if, and only if, $v = \bigvee \{u \in V \mid u \ll v \}$ for every $v \in V$; here $u \ll v$ (“$u$ totally below $v$”) means

$$\forall D \subseteq V : \quad v \leq \bigvee \quad D \implies (\exists d \in D : u \leq d).$$

Every completely distributive complete lattice in the ordinary sense is ccd, with the validity of the converse implication being equivalent to the Axiom of Choice (see [23, 24]).

\textbf{Theorem 2.8.} Let $V$ be constructively completely distributive. Then a $V$-valued closure space is a $V$-valued topological space if, and only if, its closure tower $(c^v)_{v \in V}$ satisfies

$$(C4) \quad c^\emptyset = \emptyset,$$

$$(C5) \quad c^v (A \cup B) = \cap_{u \in V} (c^u A \cup c^u B),$$

for all $v \in V, v > \bot, \text{ and } A, B \subseteq X$.

\textbf{Proof.} For the $V$-valued closure space $(X, c)$ to be $V$-valued topological means, by definition,

$$(c\emptyset)(x) = \bot \quad \text{and} \quad c(A \cup B)(x) = (cA)(x) \lor (cB)(x),$$

for all $A, B \subseteq X, x \in X$, and from Proposition 2.6 and Remark 2.7 we recall

$$(cA)(x) = \bigvee \{v \in V \mid x \in c^v A \} = \bigvee \{v \in V \mid v > \bot, x \in c^v A \}.$$

Trivially then, $(c\emptyset)(x) = \bot$ for all $x \in X$ if, and only if, $c^\emptyset = \emptyset$ for all $v > \bot$.

When $V$ is completely distributive, from (C1),(C0) one obtains, for all $v \in V, A, B \subseteq X$,

$$c^v (A \cup B) = \bigcap_{u \in V} (c^v (A \cup B)) \supseteq \bigcap_{u \in V} (c^u A \cup c^u B).$$
Hence, with the equivalences

$$\forall x \in X : \ c(A \cup B)(x) \leq (cA)(x) \lor (cB)(x)$$

$$\iff \forall x \in X, v \in V : \ (c(A \cup B)(x) \geq v \Rightarrow (cA)(x) \lor (cB)(x) \geq v)$$

$$\iff \forall x \in X, v \in V : \ (c(A \cup B)(x) \geq v \Rightarrow \forall u < v : \ (cA)(x) \geq u \lor (cB)(x) \geq u)$$

$$\iff \forall v \in V : \ c'(A \cup B) \subseteq \bigcap_{u \in V} (c'u A \cup c'u B)$$

the assertion of the Theorem follows from Proposition 2.6.

\[\square\]

**Remark 2.9.** (1) Of course, for $V$ ccd, (C0) follows from (C5) and is therefore not needed when characterizing $V$-valued topological spaces.

(2) For $v > \bot$, (C4) may be equivalently stated as

\[(C4) \quad c'0 = \bigcap_{u \in V} 0,\]

and in this form the requirement remains valid also when $v = \bot$: since there is no $u < \bot$ in $V$, trivially $\bigcap_{u \in \bot} 0 = X = c'0$ (see Remark 2.7).

When $V$ is completely distributive in the ordinary sense, then the conditions (C4), (C5) may be simplified, as follows. Recall that an element $p \in V$ is coprime if

$$\forall D \subseteq V \text{ finite : } p \leq \sqrt{D} \iff (\exists d \in D : p \leq d);$$

equivalently, if $p > \bot$, and $p \leq u \lor v$ always implies $p \leq u$ or $p \leq v$; or, equivalently, if $\{v \in V : v \nleq p\}$ is a directed subset of $V$, that is: if any of its finite subsets has an upper bound in $V$. Note that, contrary to this definition, some authors regard also $\bot$ as coprime, but that does not affect the validity of the following well-known Proposition, for which one must grant the Axiom of Choice.

**Proposition 2.10.** ([9], Theorem I-3.16) If $V$ is completely distributive, then $v = \sqrt{\{p \in V \mid p \leq v, \text{ p coprime}\}}$, for all $v \in V$.

Now we can characterize $V$-approach spaces in terms of closure towers satisfying (C1),(C2),(C3), and the following conditions (C4'), (C5'):

**Theorem 2.11.** Let $V$ be completely distributive. Then a $V$-valued closure space $(X, c)$ is a $V$-valued topological space if, and only if, its closure tower $(c^n)_{n \in V}$ satisfies

\[(C4') \quad c'0 = 0,\]

\[(C5') \quad c'(A \cup B) = c^p A \cup c^p B,\]

for all coprime elements $p \in V$ and $A, B \subseteq X$.

**Proof.** Firstly,

$$\forall x \in X : (c0)(x) = \bot \iff \forall x \in X, v \in V, v > \bot : v \nleq (c0)(x)$$

$$\iff \forall x \in X, p \in V, p \text{ coprime : } p \nleq (c0)(x) \quad \text{(Proposition 2.10)}$$

$$\iff \forall p \in V, p \text{ coprime : } c'(0) = 0.$$

Secondly, since trivially, for all coprime $p \in V, A, B \subseteq X, x \in X,$

$$x \in c^p (A \cup B) \iff p \leq c(A \cup B)(x)$$

and

$$x \in (c^p A \cup c^p B) \iff p \leq (cA)(x) \text{ or } p \leq (cB)(x)$$

$$\iff p \leq (cA)(x) \lor (cB)(x),$$
with Proposition 2.10 one obtains that
\[ \forall p \in \mathcal{V}, p \text{ coprime} : c^p(A \cup B) = c^pA \cup c^pB \iff \forall x \in X : c(A \cup B)(x) = (cA)(x) \vee (cB)(x). \]

\[ \square \]

**Example 2.12.** (1) For the terminal quantale \( 1 \) one obtains \( 1 \cdot \text{Top} = 1 \cdot \text{Cls} \cong \text{Set} \).

(2) For the two-element chain \( 2 \) (considered as a quantale with its frame structure, so that \( \varnothing = \land \)), we see that \( 2 \cdot \text{Cls} = \text{Cls} \) is the category of closure spaces, i.e., of sets \( X \) that come with an extensive, monotone and idempotent closure operation \( c : PX \rightarrow PX \), and that \( 2 \cdot \text{Top} = \text{Top} \) is the category of topological spaces (presented in terms of a finitely additive closure operation), and their continuous maps.

For Lawvere’s quantale \([0, \infty]\) we obtain from Lemma 2.5, Proposition 2.6, Remark 2.7(2) and Theorem 2.11 the following new characterizations of \([0, \infty]\)-valued closure spaces and of approach spaces.

**Corollary 2.13.** (1) A \([0, \infty]\)-valued closure space \( X \) may be described in terms of a point-set-distance function \( \delta \) satisfying

\[
\begin{align*}
\text{(R')} & \quad \delta(x, A) = 0 \text{ for all } x \in A \subseteq X, \\
\text{(T')} & \quad \delta(x, A) \leq \sup_{y \in B} \delta(y, A) + \delta(x, B) \text{ for all } A, B \subseteq X.
\end{align*}
\]

The \([0, \infty]\)-valued closure space \( X \) is a \([0, \infty]\)-valued topological space if, and only if, \( \delta \) satisfies also

\[
\begin{align*}
\text{(D)} & \quad \delta(x, \emptyset) = \infty, \\
\text{(D)} & \quad \delta(x, A \cup B) = \min\{\delta(x, A), \delta(x, B)\},
\end{align*}
\]

for all \( x \in X, A, B \subseteq X \); equivalently, if \( X \) is an approach space, so that \( \delta \) satisfies \( \text{(D1)} \)–\( \text{(D4)} \).

(2) A \([0, \infty]\)-valued closure space \( X \) is equivalently described by a closure tower \((c^\alpha : PX \rightarrow PX)_{\alpha \in [0, \infty]}\) satisfying

\[
\begin{align*}
\text{(C0)} & \quad \text{if } B \subseteq A, \text{ then } c^\alpha B \subseteq c^\alpha A, \\
\text{(C1)} & \quad \text{if } \inf_{i \in I} \beta_i \leq \alpha, \text{ then } \bigcap_{i \in I} c^{\beta_i} A \subseteq c^\alpha A, \\
\text{(C2)} & \quad A \subseteq c^\alpha A, \\
\text{(C3)} & \quad c^\alpha c^\beta A \subseteq c^{\alpha + \beta} A,
\end{align*}
\]

for all \( A \subseteq X \) and \( \alpha, \beta, \beta_i \in [0, \infty] \) (\( i \in I \)). For \( X \) to be an approach space, \((c^\alpha)_{\alpha \in [0, \infty]}\) must satisfy \( \text{(C1)} \)–\( \text{(C3)} \) and

\[
\begin{align*}
\text{(C4)} & \quad c^0(\emptyset) = \emptyset, \\
\text{(C5)} & \quad c^\alpha(A \cup B) = c^\alpha A \cup c^\alpha B,
\end{align*}
\]

for all \( A, B \subseteq X, \alpha < \infty \).

(3) A map \( f : X \rightarrow Y \) of \([0, \infty]\)-valued closure spaces \( X, Y \), presented in terms of their respective closure towers \((c^\alpha), (d^\alpha))\), is contractive if, and only if, \( f(c^\alpha A) \subseteq d^\alpha(f(A)) \) for all \( A \subseteq X, \alpha \in [0, \infty] \).

In summary, \([0, \infty]\)-\textbf{Top} = \textbf{App} is the category of approach spaces (as defined in terms of point-set-distances) that, as shown above, may be equivalently described in terms of closure towers.

The quantale \([0, \infty]\) is of course isomorphic to the unit interval \([0, 1]\), ordered by the natural \( \leq \) and provided with the ordinary multiplication as \( \otimes \). Both, \([0, \infty]\) and \([0, 1]\) are embeddable into the quantale \( \Delta \) of all distance distribution functions \( \varphi : [0, \infty] \rightarrow [0, 1] \), required to satisfy the left-continuity condition \( \varphi(\beta) = \sup_{\alpha \leq \beta} \varphi(\alpha) \), for all \( \beta \in [0, \infty] \). Its order is inherited from \([0, 1]\), and its monoid structure is given by the commutative convolution product

\[
(\varphi \circ \psi)(\gamma) = \sup_{\alpha \leq \beta \leq \gamma} \varphi(\alpha) \psi(\beta);
\]

the \( \circ \)-neutral function \( \kappa \) satisfies \( \kappa(0) = 0 \) and \( \kappa(\alpha) = 1 \) for all \( \alpha > 0 \). We note \( \kappa = \top \) in \( \Delta \) (so \( \Delta \) is integral), while the bottom element in \( \Delta \) has constant value 0; we write \( \bot = 0 \). The significance of the quantale homomorphisms \( \sigma : [0, \infty] \rightarrow \Delta \) and \( \tau : [0, 1] \rightarrow \Delta \), defined by \( \sigma(\gamma) = 0 \) if \( \gamma \leq \alpha \), and 1 otherwise, and \( \tau(u)(\gamma) = u \) if \( \gamma > 0 \), and 0 otherwise, lies in the fact that every \( \varphi \in \Delta \) has a presentation

\[
\varphi = \bigvee_{\alpha \in [0, \infty]} \sigma(\alpha) \circ \tau(\varphi(\alpha)) = \bigvee_{\alpha \in [0, \infty]} \sigma(\alpha) \circ \tau(\varphi(\alpha)).
\]
As a consequence (that was noted in [22]), one has a presentation of \( \Delta \) as a coproduct of \([0, \infty]\) and \([0, 1]\) in the category \( \text{Qnt} \) of commutative quantales and their homomorphisms, with coproduct injections \( \sigma \) and \( \tau \), respectively.

The lattice \( \Delta \) is constructively completely distributive, hence completely distributive in the presence of the Axiom of Choice. The above presentation displays \( \varphi \) as a join of coprime elements. Indeed, a distance distribution function \( \pi \) is coprime if, and only if, there are \( \alpha \in (0, \infty) \) and \( u \in [0, 1] \) such that \( \pi = \sigma(\alpha) \circ \tau(u) \), i.e., \( \pi(\gamma) = \begin{cases} 0 & \text{if } \gamma \leq \alpha, \\ u & \text{if } \gamma > \alpha. \end{cases} \)

**Remark 2.14.** Rather than basing the definition of the convolution on the multiplication in \([0, 1]\), one may consider a commutative monoid operation \( \& \) on \([0, 1]\) preserving suprema in each variable (also known as a left-continuous \( t \)-norm on \([0,1]\) and having 1 as its neutral element, such as the \( \Łukasiewicz \) operation \( \alpha \& \beta = \max\{\alpha + \beta - 1, 0\} \) or the frame operation \( \alpha \& \beta = \min\{\alpha, \beta\} \), and then define the convolution \( \odot \) on the lattice \( \Delta \) by

\[
(\varphi \odot \psi)(\gamma) = \sup_{\alpha + \beta \geq \gamma} \varphi(\alpha) \& \psi(\beta),
\]

to obtain the commutative and integral quantale \( \Delta_{k} \) (see [15]). Then \( \Delta_{k} \) is presentable as a coproduct of \([0, \infty]\) and the quantale \([0, 1]_{k} = ([0, 1], \&), 1) \), as has also been observed in [8].

**All claims that follow remain true if the convolution is being read in this more general form and \( \Delta \) being replaced by \( \Delta_{k} \).**

A probabilistic approach space \([14, 15]\) is a set \( X \) equipped with a function \( \delta : X \times PX \rightarrow \Delta \), subject to

\[
\begin{align*}
\text{(PD1)} & \quad \forall x \in X : \delta(x, \{x\}) = \kappa, \\
\text{(PD2)} & \quad \forall x \in X : \delta(x, \emptyset) = 0, \\
\text{(PD3)} & \quad \forall x \in X, A, B \subseteq X : \delta(x, A \cup B) = \delta(x, A) \lor \delta(x, B), \\
\text{(PD4)} & \quad \forall x \in X, A \subseteq X, \varphi \in \Delta : \delta(x, A) \geq (\cap_{\gamma \in \varphi} \delta(x, A)) \lor \delta(x, B), \text{ where } A^{(\varphi)} := \{x \in X \mid \delta(x, A) \geq \varphi\}.
\end{align*}
\]

Calling a map \( f : (X, \delta) \rightarrow (Y, \varepsilon) \) of probabilistic approach spaces contractive when \( \delta(x, A) \leq \varepsilon(fx, f(A)) \) for all \( x \in X, A \subseteq X \), we obtain the category \( \text{ProbApp} \).

In analogy to Corollary 2.13, the general results of this section lead to the following alternative descriptions of probabilistic approach spaces and their morphisms.

**Corollary 2.15.** (1) A function \( \delta : X \times PX \rightarrow \Delta \) is a probabilistic approach structure on a set \( X \) if, and only if, \( \delta \) satisfies (PD2),(PD3) and

\[
\begin{align*}
\text{(R*')} & \quad \delta(x, A) = \kappa \text{ for all } x \in A \subseteq X, \\
\text{(T*')} & \quad \delta(x, A) \leq (\cap_{\gamma \in \varphi} \delta(x, A)) \lor \delta(x, B) \text{ for all } A, B \subseteq X.
\end{align*}
\]

Equivalently, the function \( c : PX \rightarrow \Delta^{X} \) with \((cA)(x) = \delta(x, A)\) makes \((X, c)\) a \( \Delta \)-valued topological space.

(2) The probabilistic approach structure on a set \( X \) may be described equivalently by a family of functions \( c^{\varphi} : PX \rightarrow PX \) (\( \varphi \in \Delta \)) satisfying

\[
\begin{align*}
\text{(PC1)} & \quad \text{if } \varphi \leq \bigvee_{\gamma \in I} \psi_{i}, \text{ then } \bigcap_{i \in I} c^{\psi_{i}} A \subseteq c^{\varphi} A, \\
\text{(PC2)} & \quad A \subseteq c^{\varphi} A, \\
\text{(PC3)} & \quad c^{\varphi} c^{\psi} A \subseteq c^{\varphi \circ \psi} A, \\
\text{(PC4)} & \quad c^{\varphi}(\emptyset) = 0, \\
\text{(PC5)} & \quad c^{\varphi}(A \cup B) = c^{\varphi} A \cup c^{\varphi} B,
\end{align*}
\]

for all \( A, B \subseteq X, \varphi, \psi, \pi \in \Delta, \pi \) coprime.

(3) A map \( f : X \rightarrow Y \) of probabilistic spaces \( X, Y \), presented in terms of their respective closure towers \((c^{\varphi}), (d^{\varphi})\), is contractive if, and only if, \( f(c^{\varphi} A) \subseteq d^{\varphi}(f(A)) \) for all \( A \subseteq X, \varphi \in \Delta \).

In summary, \( \text{ProbApp} = \Delta^{\text{-Top}} \), and contractivity of a map is equivalently described by continuity with respect to \( \Delta \)-closure towers.
3. V-valued topological spaces via ultrafilter convergence

Throughout this section, the quantale \( V \) is assumed to be completely distributive.

We let \( U = (U, \Sigma, (-)) \) denote the ultrafilter monad on \( \text{Set} \). Hence, \( UX \) is the set of ultrafilters on the set \( X \), and the effect of \( U \) on a map \( f : X \rightarrow Y \) and the monad structure of \( U \) are described by

\[
Uf : UX \rightarrow UY, \quad x \mapsto f[x], \quad (B \in f[x] \iff f^{-1}B \in x),
\]

\[
(-) : X \rightarrow UX, \quad x \mapsto \check{x}, \quad (A \in \check{x} \iff x \in A),
\]

\[
\Sigma_X : UUX \rightarrow UX, \quad X \mapsto \Sigma_X, \quad (A \in \Sigma_X \iff \{x \in UX \mid A \in x\} \in X),
\]

for all \( x \in X, x \in UX, \check{x} \in UUX, A \subseteq X, B \subseteq Y \).

**Proposition 3.1.** The ultrafilter monad \( U \) distributes laxly over the \( V \)-powerset monad \( \mathbb{P}_V \), via

\[
\beta_X : U(V^X) \rightarrow V^{UX}, \quad (\beta_X \check{\varepsilon})(x) = \bigwedge \bigvee \sigma(x) \quad (\varepsilon \in U(V^X), x \in UX).
\]

**Proof.** We verify the defining conditions (a)–(f) of Section 2.

(a) With \( f[x] = \varepsilon \) one has

\[
(\beta_Y \cdot (Uf))((\beta_X \check{\varepsilon}))(\eta) = \bigwedge \bigvee \tau(y) = \bigwedge \bigvee f_i(\sigma)(fx) \geq \bigwedge \bigvee \sigma(x).
\]

Consequently,

\[
(\beta_Y \cdot (Uf))(\varepsilon)(\eta) \geq \bigwedge \bigvee \bigvee \sigma(x) = \bigwedge \bigvee \beta_X(\check{\varepsilon})(x) = ((Uf) \cdot \beta_X(\check{\varepsilon}))(\eta).
\]

(b) Since, for \( x, \eta \in UX \), one has \( y_{UX}(x)(\eta) = k \) if \( \eta = x \), and \( \perp \) otherwise, \( y_{UX} \leq \beta_X \cdot Uy_X \) follows from

\[
(\beta_X \cdot Uy_X)(x)(\eta) = \bigwedge \bigvee \sigma(x) = \bigwedge \bigvee y_X(y)(\eta) \geq k,
\]

with the last inequality following from \( A \cap B \neq \emptyset \) for all \( A, B \in x \).

(c) For all \( \omega \in U(V^X) \) and \( x \in UX \),

\[
(s_{UX} \cdot \beta_X)(\beta_V(\omega))(x) = \bigvee \Phi(\sigma) \otimes \sigma(x).
\]

Consequently, in order for us to conclude \( (s_{UX} \cdot \beta_X)(\beta_V(\omega))(x) \leq (\beta_X \cdot Uy_X)(\omega)(x) \), it suffices to show that, given any \( \varepsilon \in U(V^X), \Sigma \in \omega, A \in x, \) as well as \( u \in U\beta_V(\omega), \nu \in \beta_X(\check{\varepsilon})(x) \) in \( V \), that there are \( \Phi \in S, \sigma \in V^X, x \in A \) with \( u \otimes \nu \leq \Phi(\sigma) \otimes \sigma(x) \). But indeed, from the stated hypothesis on \( u, \nu \in V \), for all \( \Sigma \in \varepsilon \) one obtains \( \Phi_\Sigma \in S, \sigma_\Sigma \in S \) with \( u \leq \Phi_\Sigma(\sigma_\Sigma) \), and \( \tau_\Sigma \in S, x_\Sigma \in A \) with \( \nu \leq \tau_\Sigma(x_\Sigma) \). Now, the set \( M = \{\sigma_\Sigma \mid \Sigma \in \varepsilon\} \) satisfies \( M \cap S ) \neq \emptyset \) for all
$S \in s$ and must therefore belong to $s$ (since, otherwise, we could find an ultrafilter properly containing $s$); likewise, $N = \{ \tau_S \mid S \in s \} \in s$. Consequently, $M \cap N \neq \emptyset$, from which one derives the needed claim.

(d) Let $\sigma \in \mathcal{V}^X$, $x \in \mathcal{U}X$. If $x = \cdot$ for $x \in X$, then

$$(-)_1(\sigma)(x) = \bigvee_{y \in y(x)} \sigma(y) = \sigma(x) = \bigwedge_{x \in x} \tau(y) = \beta_X(\cdot)(x);$$

otherwise $(-)_1(\sigma)(x) = \bot$, and the needed inequality holds trivially.

(e) If for $X \in \mathcal{U}UUX$, $\Xi \in \mathcal{U}U(UV^X)$ we have $\Sigma_X(X) = x$, $\Sigma_X(\Xi) = s$, then for any given $S \in s$, $A \in x$, there are $S_0 \in \Xi$, $\mathcal{A}_0 \in X$ such that $S \in t$, $A \in \eta$ for all $t \in S_0$, $\eta \in \mathcal{A}_0$. Obviously then,

$$\bigvee_{x \in x} \bigwedge_{y \in y(x)} \tau(y) \leq \bigwedge \sigma(x).$$

Consequently, for all $\Xi \in \mathcal{U}U(UV^X)$, $x \in \mathcal{U}X$, putting $s = \Sigma_X(\Xi)$ one obtains

$$((\Sigma_X) \cdot \beta_{\mathcal{U}X} \cdot \mathcal{U}H_X(\Xi))(x) = \bigvee_{x \in x} \bigwedge \varphi(y) = \bigvee_{S \in \mathcal{S}_X} \bigwedge_{S \in \mathcal{S}_X} \bigwedge_{S \in \mathcal{S}_X} \tau(y) \leq \bigwedge \sigma(x) = \beta_X \cdot \Sigma_X(\Xi)(s).$$

(f) For $g, h : Z \to \mathcal{V}^X$ with $g \leq h$ and all $x \in \mathcal{U}Z$, $x \in \mathcal{U}X$, one has

$$(\beta_X \cdot \mathcal{U}g(\xi))(x) = \bigvee_{x \in x} (g \xi)(x) \leq \bigwedge_{x \in x} (h \xi)(x) = \beta_X \cdot \mathcal{U}h(\xi)(x).$$

\[\square\]

**Remark 3.2.** The lax extension $\mathcal{U} : \mathcal{V}^X \to \mathcal{V}^X$ of $\mathcal{U}$ corresponding to $\beta$ (see Remark 2.1) is given by

$$\mathcal{U}r(x, \eta) = \bigwedge_{A \in A} \bigwedge_{B \in B} r(x, y),$$

for all $r : X \to Y$, $x \in \mathcal{U}X$, $\eta \in \mathcal{U}Y$.

A straightforward calculation gives a description of lax $(\beta, \mathcal{V})$-algebras (with $\beta$ as in Proposition 3.1), in analogy to the description of $(\alpha, \mathcal{V})$-algebras of Lemma 2.5.

**Lemma 3.3.** A map $\ell : \mathcal{U}X \to \mathcal{V}^X$ makes $(X, \ell)$ a lax $(\beta, \mathcal{V})$-algebra if and only if, $\ell$ satisfies

(R*) $\forall x \in X : \ell^k(x) \leq \ell(x)(x),$

(T*) $\forall x \in \mathcal{U}X, \eta \in \mathcal{U}X, z \in X : \bigwedge_{A \in A \cap X} \bigwedge_{B \in B \cap X} (\ell x)(y) \otimes (\ell y)(z) \leq \ell(x,y)(z).$
A map \( f : X \to Y \) is a lax homomorphism \( f : (X, \ell) \to (Y, \ell') \) of lax \((\beta, \mathcal{V})\)-algebras if, and only if,

\[ \forall x \in UX, y \in X : \quad (\ell)(x)(y) \leq (\ell' f(x))(f(y)). \]

When displayed in terms of the \( \mathcal{V} \)-relation \( a : UX \to X \) with \( a(x, y) = (\ell)(x)(y) \), (see Remark 2.4), conditions \((R^*)\), \((T^*)\) read as

\[ k \leq a(\Delta(x, y)) \quad \text{and} \quad \overline{a}(\Delta(x, y)) \otimes a(\Delta(y, z)) \leq a(\Delta(x, z)), \quad (*) \]

for all \( x, z \in X, \eta \in UX, \xi \in UX \).

Next we will establish an adjunction between the categories \((\alpha, \mathcal{V})\)-\textbf{Alg} \( \equiv (\mathcal{P}, \mathcal{V}, \mathcal{P})\)-\textbf{Cat} \( \equiv \mathcal{V}\text{-Cls} \) and \((\beta, \mathcal{V})\)-\textbf{Alg} \( \equiv (\mathcal{U}, \mathcal{V}, \mathcal{U})\)-\textbf{Cat}, the restriction of which will then give an isomorphism

\[ \mathcal{V}\text{-Top} \equiv (\beta, \mathcal{V})\text{-Alg}. \]

First recall from [22] that, given lax extensions \( \hat{S}, \hat{T} \) of \textbf{Set}-monads \( S, T \) to \textbf{V-Rel}, an algebraic morphism \( h : (\hat{S}, \hat{T}) \to (S, T) \) is a family of \( \mathcal{V} \)-relations \( h_X : SX \to TX \) \((X \in \text{Set})\), satisfying the following conditions for all \( f : X \to Y \) in \text{Set} and \( r : X \to Y, a : TX \to X \) in \text{V-Rel}:

\begin{align*}
& a. \quad Tf \circ h_X \leq h_Y \circ Sf, \quad \text{(lax naturality)} \\
& b. \quad e_X \leq h_X \circ d_X, \quad \text{(lax unit law)} \\
& c. \quad m_X \circ h_X \circ \hat{S}h_X \leq h_Y \circ n_X, \quad \text{(lax multiplication law)} \\
& d. \quad h_Y \circ \hat{T}r \leq \hat{T}r \circ h_X, \quad \text{(lax compatibility)} \\
& e. \quad \hat{S}(a \circ h_X) \leq \hat{S}a \circ \hat{S}h_X, \quad \text{(strictness of } \hat{S} \text{ at } h) \\
\end{align*}

Here, for \( s : Y \to Z \), the composite \( s \circ r : X \to Y \) in \text{V-Rel} is given by \((s \circ r)(x, z) = \bigvee_{y \in Y} s(y, z) \otimes r(x, y)\), and (as in [12]) we identify a map \( f : X \to Y \) with its \( \mathcal{V} \)-graph \( f_* : X \to Y \), given by \( f_*(x, y) = k \) if \( f(x) = y \), and \( \perp \) otherwise.

Now, such lax transformation \( h : (\hat{S}, \hat{T}) \to (S, T) \) induces the algebraic functor

\[ A_h : (\mathcal{V}, \mathcal{V})\text{-Cat} \to (\mathcal{S}, \mathcal{S})\text{-Cat}, \quad (X, a) \mapsto (X, a \circ h_X). \]

Considering \( S = \mathcal{P}, T = \mathcal{U} \), let us consider \( e_X : PX \to UX \) by

\[ e_X(A, x) = \begin{cases} 
\kappa & \text{if } A \in x \\
\perp & \text{otherwise}
\end{cases}, \]

for all \( A \subseteq X, x \in UX \).

**Proposition 3.4.** \( e : (\mathcal{P}, \hat{\mathcal{P}}) \to (\mathcal{U}, \hat{\mathcal{U}}) \) is an algebraic morphism and, hence, induces the algebraic functor

\[ A_e : (\mathcal{U}, \mathcal{V})\text{-Cat} \to (\mathcal{P}, \mathcal{V}, \hat{\mathcal{P}})\text{-Cat}, \quad (X, a) \mapsto (X, a \circ e_X). \]

**Proof.** We verify conditions a–e above,

a. Trivially, if \( A \subseteq X \subseteq UX \) and \( f[x] = \eta \), then \( f(A) \subseteq \eta \), and \((Uf \circ e_X)(A, \eta) \leq (e_Y \circ P_f)(A, \eta)\) follows.

b. Likewise, if \( x = \chi \), then \( [x] \subseteq x \), and \((\cdot)(x, x) \leq e_X \circ [\cdot](x, x)\) follows.

c. For \( \overline{x} \in \mathcal{P}X, A \subseteq \mathcal{P}X \) one has

\[ (e_{\mathcal{P}X} \circ \hat{\mathcal{P}}e_X)(A, \overline{x}) = \bigvee_{x \in \mathcal{P}X} \hat{\mathcal{P}}e_X(A, B) \otimes e_{\mathcal{P}X}(B, \overline{x}) = \bigvee_{x \in \mathcal{P}X} \hat{\mathcal{P}}e_X(A, B) = \bigvee_{x \in X} \bigvee_{y \in Y} e(A, y) = k \]

in the case that, for all \( \eta \in \mathcal{B} \), there is \( A \in \mathcal{A} \) with \( A \subseteq \eta \), and \( \perp \) otherwise. So, in the former case, given any \( B \in X \), one has \( B \subseteq \{ \eta \} \subseteq \mathcal{P}X \mid \bigcup_{A \in \eta} \subseteq \overline{x} \) and, hence, \( \bigcup_{x \in \mathcal{P}X} \subseteq \Delta X \). Consequently, \((\Delta X \circ e_{\mathcal{P}X} \circ \hat{\mathcal{P}}e_X)(A, x) \leq (e_X \circ \Delta X)(A, x)\).

d. For \( r : X \to Y, A \subseteq X, \eta \subseteq UX \), we must compare

\[ (e_Y \circ \hat{\mathcal{P}}r)(A, \eta) = \bigvee_{B \subseteq Y} e_Y(B, \eta) \otimes \hat{\mathcal{P}}r(A, B) = \bigvee_{B \subseteq Y} \bigvee_{\eta \subseteq B} \bigvee_{a \in A} r(x, y) \]
with 
\[(\bar{U}r \circ e_X)(A, \eta) = \bigvee_{x \in UX} \bar{U}r(x, \eta) \otimes e_X(A, x) = \bigvee_{\eta \in \eta A, y \in X} \bigvee_{x \in X} r(x', y') .\]

So, given \(B \in \eta\), we consider \(u \leq \bigwedge_{\eta \in \eta A} r(x, y)\) in \(V\). For all \(y \in B\) we may then pick \(fy \in A\) with \(u \leq r(fy, y)\). With the map \(f : B \to A\) we choose an ultrafilter \(x\) on \(X\) that contains all the sets \(f(C)\), \(C \in \eta, C \subseteq B\). For all such \(C\) and any \(B' \in \eta\), since \(C \cap B' \neq \emptyset\), we finally obtain some \(y' \in B'\) with \(u \leq r(fy', y')\). Now \((a \circ \hat{u})r(B, \eta) \leq (\bar{U}r \circ e_X)(A, \eta)\) follows.

e. For \(A \subseteq PX, B \subseteq X\) one has 
\[
\hat{P}(a \circ e_X)(A, B) = \bigvee_{y \in B} \bigwedge_{\alpha \in A} a(x, y) \otimes e_X(A, x) = \bigvee_{\eta \in \eta A} \bigvee_{y \in B} a(x, y),
\]
while 
\[
(\hat{P}a \circ \hat{P}e_X)(A, B) = \bigvee_{y \in B} \hat{P}a(y, B) \otimes \hat{P}e_X(A, B) = \bigvee_{\eta \in \eta A} \bigwedge_{y \in B} a(x, y) \otimes \bigwedge_{y \in B} e_X(A, y').
\]

Now, whenever \(u \leq (\hat{P}a \circ \hat{P}e_X)(A, B)\) in \(V\), for all \(y \in B\) one obtains \(x_y \in A_y \in A\) with \(u \leq a(x_y, y)\). Putting \(B := \{x, \mid y \in B\}\) one sees \(u \leq (\hat{P}a \circ \hat{P}e_X)(A, B)\), which gives the needed inequality.

\[\square\]

**Remark 3.5.** In the quantaloid (see [21]) \(V\)-Rel, the sup-map 
\[
V\text{-Rel}(UX, X) \to V\text{-Rel}(PX, X), \quad a \mapsto a \circ e_X,
\]
has a right adjoint, which assigns to \(\delta : PX \to X\) the \(V\)-relation \(\delta \circ e_X : UX \to X\), given by 
\[
(\delta \circ e_X)(t, x) = \bigwedge_{A \in X} (\delta(A, x) \circ e_X(A, x)) = \bigwedge_{A \in X} \delta(A, x).
\]

Writing \((ca)(x)\) for \(\delta(A, x)\) we will take advantage of this obvious fact in the proof of the Theorem below.

If we describe \((P, V)\)-categories as \(V\)-valued closure spaces and \((U, V)\)-categories as lax \((\beta, V)\)-algebras then, \(A_x\) takes the form 
\[
A_x : (\beta, V)\text{-Alg} \to V\text{-Cls}, \quad (X, \ell) \mapsto (X, c_\ell : PX \to V^X), \quad (c_\ell(A)(x) = \bigvee_{x \in UX, x \in A} (f x)(x).
\]

We are now ready to prove the main result of this paper:

**Theorem 3.6.** For a completely distributive quantale \(V\), the algebraic functor \(A_x\) embeds \((\beta, V)\)-Alg into \(V\)-Cls as a full coreflective subcategory, which is precisely the category \(V\text{-Top}^v\) of \(V\)-valued topological spaces.

**Proof.** That \(A_x\) actually takes values in \(V\)-\textbf{Top} is just a reflection of the fact that, for any ultrafilter \(x\) on a set \(X\), one has \(\emptyset \notin x\), and \(A \cup B \in x\) only if \(A \in x\) or \(B \in x\). Next we prove that \(A_x\) has a right adjoint, described by (see Remark 3.5) 
\[
R : V\text{-Cls} \to (\beta, V)\text{-Alg}, \quad (X, c) \mapsto (X, \ell_c : UX \to V^X), \quad (\ell_c(x)(x) = \bigwedge_{A \in X} (cA)(x).
\]

Given \((X, c) \in V\text{-Cls}\) we must first show \((X, \ell_c) \in (\beta, V)\text{-Alg}\), that is: writing \(a(x, y)\) for \((\ell_c(x)(y)\), we must establish \((*)\) of Lemma 3.3. Trivially, \(k \leq \bigwedge_{A \in X} (cA)(x) = a(x, x)\) for all \(x \in X\). Our strategy to show 
\[
\bar{U}a(x, \eta) \otimes a(\eta, z) \leq a(\Sigma x, z) = \bigwedge_{A \in X} (cA)(z)
\]
for all \(x \in UX, \eta \in UX, z \in X\), is to consider any \(A \in \Sigma \tilde{x}\) and 
\[
u \leq \bar{U}a(x, \eta) = \bigwedge_{\tilde{x} \in X} \bigvee_{x \in X} a(x, y)
\]
in $\mathcal{V}$ and to show $u \otimes a(\eta, z) \leq (cA)(z)$. Indeed, with $A \in \Sigma X$ one has $\mathcal{A} := \{x \in UX \mid A \in x\} \in \mathfrak{X}$. Then, putting $C := \{y \in X \mid \exists x \in \mathcal{A} : u \leq a(x, y)\}$, since $a(x, y) \leq (cA)(y)$ whenever $A \in x$, we obtain $C \subseteq c^u A = \{y \in X \mid (cA)(y) \geq u\}$. Since $u \ll \bigwedge_{B \in \mathcal{B}} \vee_{x \in \mathcal{A} \cap cB} a(x, y)$, so that for all $B \in \eta$ there is $x \in \mathcal{A}$ with $u \leq a(x, y)$, we see that $B \cap C \neq \emptyset$ whenever $B \in \eta$. Maximality of $\eta$ therefore forces $C \in \eta$, and then $c^u A \in \eta$. With (C3) of Remark 2.7 we conclude

$$u \otimes a(\eta, z) = u \otimes \bigwedge_{B \in \mathcal{B}} (cB)(z) \leq u \otimes c(c^u A)(z) \leq (cA)(z).$$

For the adjunction $A_c + R$, it now suffices to show that, given a $\mathcal{V}$-valued closure space $(X, c)$ and a $(\beta, \mathcal{V})$-algebra $(Y, f)$, a map $f : Y \rightarrow X$ is a morphism $A_2(Y, \ell) \rightarrow (X, c)$ in $\mathcal{V}$-Cls if, and only if, it is a morphism $(Y, f) \rightarrow R(X, c)$ in $(\beta, \mathcal{V})$-Alg. This, however, is obvious, since either statement means equivalently

$$\forall B \in \eta \in \mathcal{U}Y, y \in Y : (\ell \eta)(y) \leq c(f(B))(fy).$$

Next, for a $\mathcal{V}$-valued topological space $(X, c)$, we must show $A_2 R(X, c) = (X, c)$, that is: $c_{\ell_2} = c$. Since the adjunction gives $c_{\ell_2} \leq c$, it suffices to show “$\geq$”, that is: for all $A \subseteq X, x \in X$,

$$(cA)(x) \leq \bigvee_{x \in \mathcal{A}} (cB)(x),$$

and for that, by Proposition 2.10, it suffices to check that every coprime element $p$ in $\mathcal{V}$ with $p \leq (cA)(x)$ satisfies $p \leq \bigvee_{x \in \mathcal{A}} \bigwedge_{B \in \mathcal{B}_p} (cB)(x)$. But the set $I_p = \{B \subseteq X : (\ell B, x) \neq \emptyset\} \subseteq PX$ is directed since $p$ is coprime, and $I_p$ is disjoint from the filter $\{B \subseteq X : A \subseteq B\}$. There is therefore an ultrafilter $x_p$ with $A \in x_p$ disjoint from $I_p$. Thus, for all $B \in x_p$, $(cB)(x) \geq p$ and, consequently,

$$p \leq \bigwedge_{B \in x_p} (cB)(x) \leq \bigvee_{x \in \mathcal{A}} \bigwedge_{B \in \mathcal{B}_p} (cB)(x).$$

Finally we show $RA_2(X, \ell) = (X, \ell)$ for every $(X, \ell) \in (\beta, \mathcal{V})$-Alg, that is: $\ell_{c_2} = \ell$. As the adjunction gives “$\geq$”, we need to show only $\ell_{c_2} \leq \ell$. Writing $a(x, y)$ for $(\ell \eta)(y)$, this means that, for all $x \in UX, x \in X$, we must prove

$$(\ell_{c_2} x)(x) = \bigwedge_{A \in \mathcal{A}} \bigvee_{x \in \mathcal{A}} a(\eta, x) \leq a(x, x).$$

To this end, considering any $u \ll (\ell_{c_2} x)(x)$ in $\mathcal{V}$, for all $A \in x$ one obtains $\eta_A \in UX$ with $A \in \eta_A$ and $u \leq a(\eta_A, x)$. So, for all $A \in x$, the sets

$$\mathcal{A}_A = \{\eta \in UX \mid A \in \eta, u \leq a(\eta, x)\}$$

are not empty, and we can choose an ultrafilter $\mathfrak{X}$ on $UX$ containing all of them. Since for every $A \in x$ one has $\{\eta \in UX \mid A \in \eta\} \supseteq \mathcal{A}_A \in \mathfrak{X}$, we obtain $\Sigma \mathfrak{X} = x$. Furthermore,

$$\bigvee_{A \in \mathcal{A}} \bigwedge_{B \in \mathfrak{X}} a(\eta, y) = \bigwedge_{A \in \mathcal{A}} \bigvee_{x \in \mathfrak{X}} a(\eta, x) = \bigwedge_{A \in \mathcal{A}} \bigvee_{x \in \mathfrak{X}} a(y, x) \geq u.$$

With the transitivity of $a$ we conclude

$$a(x, x) = a(\Sigma \mathfrak{X}, x) \geq \bigvee_{A \in \mathcal{A}} a(x, x) \geq u \otimes k = u,$$

and $a(x, x) \geq (\ell_{c_2} x)(x)$ follows, as desired. \qed

The isomorphism

$$\mathcal{V}$-$\text{Top} \equiv (\mathcal{U}, \mathcal{V}, \mathcal{U})$-$\text{Cat} \equiv (\beta, \mathcal{V})$-$\text{Alg}$$

gives Barr’s [1] description of topological spaces and the Clementino-Hofmann [4] presentation of approach spaces in terms of ultrafilter convergence when one chooses $\mathcal{V} = 2$ and $\mathcal{V} = [0, \infty]$, respectively. For $\mathcal{V} = \Delta$ we obtain the corresponding description of ProbApp, as follows.
Corollary 3.7. The structure of a probabilistic approach space on a set $X$ may be described equivalently as a map $\ell : UX \rightarrow AX$ satisfying, for all $x \in UX, y \in UX, z \in X$,

\begin{align*}
(R^*)\; &\kappa \leq (\ell x)(x), \\
(T^*)\; &\left( \bigwedge_{\alpha \leq \beta} \bigvee_{\gamma \leq \delta} (\ell x)(y) \right) \circ (\ell y)(z) \leq \ell (\Sigma X)(z).
\end{align*}

A map $f : X \rightarrow Y$ of probabilistic approach spaces $(X, \ell), (Y, \ell')$ is contractive precisely when, for all $x \in UX, x \in X,$

\begin{align*}
(M^*)\; & (\ell x)(x) \leq (\ell' f x)(f x).
\end{align*}

4. Change-of-base functors

For a monad $T = (T, m, c)$ and a quantale $V$, let us call a set $X$ equipped with a map $c : TX \rightarrow VX$ a $(T, V)$-graph. With a morphism $f : (X, c) \rightarrow (Y, d)$ required to satisfy (M) of Definition 2.3, we obtain the category $(T, V)$-Gph, which contains $(\lambda, V)$-Alg as a full subcategory, for any lax distributive law $\lambda$ of $T$ over $P_V$. For a monotone map $\varphi : V \rightarrow W$ one has the change-of-base functor

$$B_\varphi : (T, V)$-Gph \rightarrow (T, W)$-Gph, \quad (X, c) \mapsto (X, \varphi^X \cdot c),$$

with $\varphi^X : VX \rightarrow WX, \sigma \mapsto \varphi \cdot \sigma$. Since $(T, V)$-graphs actually refer neither to the monad structure of the functor $T$ nor to the quantalic structure of the lattice $V$, they behave well under any adjunction of monotone maps:

Lemma 4.1. If $\varphi \dashv \psi : W \rightarrow V$, then $B_\varphi \dashv B_\psi : (T, W)$-Gph \rightarrow (T, V)$-Gph.

Proof. Given a $(\mathbb{T}, V)$-graph $(X, c)$ and a $(\mathbb{T}, W)$-graph, we must verify that a map $f : X \rightarrow Y$ is a morphism $(X, c) \rightarrow (Y, \psi^V \cdot d)$ if, and only if, it is a morphism $(X, \varphi^X \cdot c) \rightarrow (Y, d)$, which amounts to showing

$$f_1(W) \cdot \varphi^X \cdot c \leq d \cdot T f \iff f_1(W) \cdot \varphi^X \cdot c \leq d \cdot f_1(V).$$

But this is obvious: given the left-hand inequality, compose it from the left with $\varphi^Y$ and use $\varphi^Y \cdot \psi^Y \leq 1_{\mathbb{P}^V}$ and $f_1(W) \cdot \varphi^X \leq \varphi^Y \cdot f_1(V)$ to obtain the right-hand inequality. The converse direction is similar. \hspace{1cm} \Box

If we are given lax distributive laws $\lambda, \kappa$ of $T$ over $P_V, P_W$, respectively, what it takes for $B_\varphi$ to map $(\lambda, V)$-Alg into $(\kappa, W)$-Alg is well known from the context of $(\mathbb{T}, V)$-categories (see [12, 22]): $\varphi : (V, \otimes, k) \rightarrow (W, \otimes, l)$ should be a lax homomorphism of quantales, that is: monotone, with $l \leq k$ and $\varphi(u \otimes v) \leq \varphi(u) \otimes \varphi(v)$, for all $u, v \in V$; in addition, $\varphi$ should satisfy the $\lambda$-$\kappa$-compatibility condition $\varphi^X \cdot \lambda_X \geq \kappa_Y \cdot T(\varphi^X)$.

However, in order to be able to restrict the adjunction of Lemma 4.1 to the categories of lax algebras, while $\psi$ needs to satisfy these conditions, one does not need to require any additional condition (beyond monotonicity) on its left adjoint $\varphi$, thanks to the following simple fact:

Lemma 4.2. For a lax distributive law $\lambda$ of $T$ over $P_V, (\lambda, V)$-Alg is reflective in $(\mathbb{T}, V)$-Gph. The reflector assigns to a $(\mathbb{T}, V)$-graph $(X, c)$ the lax $(\lambda, V)$-algebra $(X, \overline{c})$, with

$$\overline{c} = \bigwedge \{ c' : TX \rightarrow VX \mid c \leq c', (X, c') \in (\lambda, V)$-Alg\}.$$

Proof. As infima in $V^X$ are formed pointwise, and as $\lambda$ is monotone, when all $c' \geq c$ satisfy (R), (T), the same is true for $\overline{c}$, since

$$c' \cdot e_X \geq y_X \quad \text{and} \quad c' \cdot m_X \geq s_X \cdot c' \cdot \lambda_X \geq s_X \cdot (\overline{c}) \cdot \lambda_X.$$ 

Furthermore, for any morphism $f : (X, c) \rightarrow (Y, d)$ one has $c \leq f^1 \cdot d \cdot T f$, where $f_1 \cdot f^1 : V^Y \rightarrow VX$. Since $c \leq f^1 \cdot d \cdot T f$ satisfies (R), (T) when $d$ does, in that case one has $\overline{c} \leq f^1 \cdot d \cdot T f$, and therefore a morphism $f : (X, \overline{c}) \rightarrow (Y, d)$. \hspace{1cm} \Box

Proposition 4.3. For lax distributive laws $\lambda, \kappa$ of a monad $T$ over $P_V, P_W$, respectively, and a lax homomorphism $\psi : W \rightarrow V$ that preserves infima and satisfies the $\lambda$-$\kappa$-compatibility condition, the change-of-base functor

$$B_\psi : (\kappa, W)$-Alg \rightarrow (\lambda, V)$-Alg

has a left adjoint $\overline{B}_\psi$, given by $(X, c) \mapsto (X, \overline{c} \cdot c)$, where $\varphi \dashv \psi$. \hspace{1cm} \Box
Proof. As an infima-preserving map of complete lattices, ψ does indeed have left adjoint ϕ. Following Lemma 4.1 and 4.2, the left adjoint of the functor \( B_ϕ \) is just the composite of the two left adjoints established previously. \( \square \)

We can now apply the Proposition to the ultrafilter law \( \beta = \beta^{(V)} \) of Proposition 3.1, first noting that any map \( \varphi : V \rightarrow W \) satisfies the \( \beta^{(V)}, \beta^{(W)} \)-compatibility condition—strictly so, as a quick inspection reveals. As usual, we write \( V-Cat \) for \( ( \mathbb{L}, V )-Cat \equiv ( 1, V )-Alg \) (where \( \mathbb{L} \) is the identical monad on \( Set \) and \( 1 : P_V \rightarrow P_V \) the identical transformation). Recall that a quantale \( V \) is integral when its \( \otimes \)-neutral element \( k \) is the top element \( \top \) in \( V \).

**Theorem 4.4.** For completely distributive and integral quantales \( V, W \), let \( \varphi : V \rightarrow W \) be monotone and \( \psi : W \rightarrow V \) a lax homomorphism of quantales. Then the following statements are equivalent:

(i) \( \varphi \cdot \psi : W \rightarrow V \);
(ii) \( \overline{B}_ϕ \cdot B_ϕ : ( β, W )-Alg \rightarrow ( β, V )-Alg \);
(iii) \( \overline{B}_ϕ \cdot B_ϕ : V-Cat \rightarrow V-Cat \).

Proof. The implications (i)⇒(ii) and (i)⇒(iii) follow from Proposition 4.3.

For (iii)⇒(i) we must show \( v \leq \psi \varphi(v) \) and \( \varphi \psi(w) \leq w \), for all \( v \in V, w \in W \). Consider \( X = \{ x, y \} \) with \( x \neq y \) and, for any \( v \in V \), define a \( V \)-category structure \( a : X \rightarrow X \) on \( X \) by \( a(x, x) = a(y, y) = k \) and \( a(x, y) = a(y, x) = v \). Since \( W \) is integral, one easily sees that the least \( W \)-category structure \( b \) on \( X \) with \( \varphi a \leq b \) is given by \( b(x, x) = b(y, y) = \top \) and \( b(x, y) = b(y, x) = \varphi(v) \); hence, \( \overline{B}_ϕ(X, a) = ( X, b ) \). Since the adjunction unit \( ( X, a ) \rightarrow B_ϕ \overline{B}_ϕ(X, a) = ( X, \psi b ) \) is a \( V \)-functor, \( v = a(x, y) \leq \psi b(x, y) = \psi \varphi(v) \) follows. Similarly one shows \( \varphi \psi(w) \leq w \), and (i) follows.

For (ii)⇒(i), one may proceed as in (iii)⇒(i), simply because, for finite \( X \), one has \( UX \equiv X \). \( \square \)

We note that the equivalence of (i) and (ii) appears in [7], Theorem 3.1, under the hypothesis that both \( \varphi \) and \( \psi \) be lax homomorphisms of quantales.

For the sake of completeness we also note that the hypothesis of Theorem 4.4 that \( \psi \) be a lax homomorphism, comes for free when \( \varphi \) is a homomorphism of quantales, i.e., a sup-preserving map which also preserves the monoid structure of the quantales, thanks to the following proposition.

**Proposition 4.5.** When \( W \) is integral, the right adjoint of a homomorphism \( \varphi : V \rightarrow W \) of quantales is a lax homomorphism of quantales.

Proof. Let \( \psi : W \rightarrow V \) be the right adjoint of the sup-preserving map \( \varphi \). Since \( \psi \) preserves infima, \( \psi(\top) = \top \geq k \).

Also, for all \( u, w \in W \),

\[
\psi(u \otimes w) \geq \psi(\varphi \psi(u) \otimes \varphi \psi(w)) = \psi \varphi(\psi(u) \otimes \psi(w)) \geq \psi(u) \otimes \psi(w).
\]

\( \square \)

We will now apply Proposition 4.3 and Theorem 4.4 to the only homomorphism \( \iota : 2 = \{ \bot, \top \} \rightarrow V \), given by \( \iota(\bot) = \bot, \iota(\top) = k \). The monotone map \( \iota \) has a right adjoint \( \pi \), given by \( ( \pi v ) = \top \iff v \geq k \) for all \( v \in V \), which is a lax homomorphism of quantales. If \( V \) is integral, \( \iota \) has also a left adjoint \( \alpha \), given by \( ( \alpha v ) = \bot \iff v = \bot \) for all \( v \in V \). Considering the identical monad \( \mathbb{I} \) one obtains the well-known fact (see [12]) that, for \( V \) non-trivial, \( B_\iota \) embeds the category \( \text{Ord} \) of (pre)ordered sets and monotone maps as a full coreflective subcategory into \( V-Cat \), which is also reflective when \( V \) is integral. Complete distributivity of \( V \) is not needed for this, but it becomes essential now when we consider the ultrafilter monad \( U \) and its lax distributive law \( \beta \).

**Corollary 4.6.** For a non-trivial completely distributive quantale \( V \), \( B_\iota \) embeds the category \( \text{Top} \) of topological spaces into \( V-\text{Top} \) as a full coreflective subcategory, with coreflector \( B_\alpha \). If \( V \) is integral, the embedding is also reflective, with reflector \( \overline{B}_\iota \). In particular, \( \text{Top} \) is both, reflective and coreflective, in \( \text{App} \), as well as in \( \text{ProbApp} \).

In fact, using the same technique as above we can refine the last statement of the Corollary and show:

**Corollary 4.7.** \( \text{App} \) is fully embedded into \( \text{ProbApp} \) as a reflective and coreflective subcategory.
Proof. The homomorphism $\sigma : [0, \infty] \to \Delta$ defined in Section 2 (after Corollary 2.13), has a right adjoint

$$\rho : \Delta \to [0, \infty], \varphi \mapsto \inf\{\alpha \in [0, \infty] | \varphi(\alpha) = 1\},$$

which is a lax homomorphism of quantales, as well as a left adjoint

$$\lambda : \Delta \to [0, \infty], \varphi \mapsto \sup\{\alpha \in [0, \infty] | \varphi(\alpha) = 0\}.$$ 

One therefore has the adjunctions

$$\overline{B}_\rho \dashv B_\sigma \dashv B_\rho : (\Delta, \Alg) \to (\beta, [0, \infty], \Alg).$$

\[ \square \]

Remark 4.8. As already stated in general in Remark 2.14, the proof of Corollary 4.7 remains valid if we equip the set of distance distribution functions with the monoidal structure

$$(\varphi \otimes \psi)(\gamma) = \sup_{\alpha \leq \gamma} \varphi(\alpha) \& \psi(\beta),$$

where $\&$ is any left-continuous continuous $t$-norm on $[0, 1]$. In the case $\& = \land$, the corresponding proof was first carried out by Jäger [15].

Here is a third application of Theorem 4.4:

Example 4.9. Let $\text{DnV}$ be the set of all down-closed subsets of $V$ which, when ordered by inclusion, is a completely distributive lattice. It becomes a quantale with

$$A \odot B = \{c \in V | \exists a \in A, b \in B : c \leq a \odot b\} \quad (A, B \in \text{DnV})$$

and $\odot$-neutral element the down-closure $\downarrow k$ of the $\odot$-neutral element $k$ of $V$. $\text{DnV}$ is integral if, and only if, $V$ is integral. More importantly, if $V$ is completely distributive, we have adjunctions

$$\downarrow \dashv \text{sup} \dashv \downarrow : V \to \text{DnV}$$

(see [23]). Furthermore, $\text{sup} : \text{DnV} \to V$ is a homomorphism of quantales, while $\downarrow$ is a lax homomorphism (but never a homomorphism, as it fails to preserve the bottom element). Therefore, we obtain the adjunctions

$$\overline{B}_\delta \dashv B_{\text{sup}} \dashv B_\delta : \text{V-Cat} \to \text{DnV-Cat} \quad \text{and} \quad \overline{B}_\delta \dashv B_{\text{sup}} \dashv B_\delta : (\beta, \text{V}, \Alg) \to (\beta, \text{DnV}, \Alg).$$

For $V$ an $n$-element chain, $\text{DnV}$ is an $(n + 1)$-element chain, which contains two distinct copies of $V$, one reflectively embedded, the other coreflectively. If $n > 2$, $V$-categories are generalized (pre)ordered sets $X$, for which the truth value for two points in $X$ being related allows for a discrete linear range, beyond $\top$ or $\bot$. For $V = ([0, \infty], \geq)$, in addition to the order embedding

$$\downarrow : [0, \infty] \to \text{Dn}[0, \infty], \alpha \mapsto [\alpha, \infty],$$

which preserves the monoidal structure, but is not a sup-map, one has the order embedding

$$\forall : [0, \infty] \to \text{Dn}[0, \infty], \alpha \mapsto (\alpha, \infty],$$

which is a sup-map, but does not preserve the monoidal structure. Since $\text{Dn}[0, \infty]$ is a disjoint union of the images of the two order embeddings, a $\text{Dn}[0, \infty]$-category structure on a set $X$ will return to a pair of points in $X$ one of two types of distances, with one type always ranking below the other, despite having equal numerical value (since $(\alpha, \infty] \subseteq [\alpha, \infty]$, for all $\alpha \in [0, \infty]$).

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